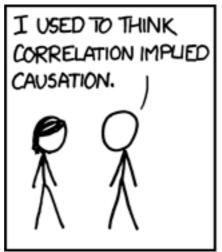


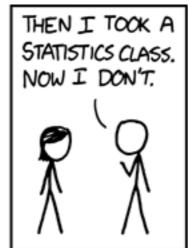
#### Scalable Machine Learning

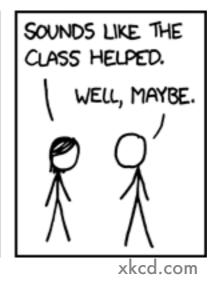
2. Statistics

#### Alex Smola Yahoo! Research and ANU

http://alex.smola.org/teaching/berkeley2012 Stat 260 SP 12





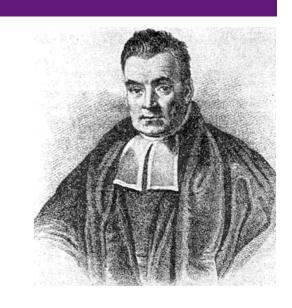


2. Statistics

Essential tools for data analysis

#### Statistics

- Probabilities
  - Bayes rule, Dependence, independence, conditional probabilities
  - Priors, Naive Bayes classifier
  - Tail bounds
    - Chernoff, Hoeffding, Chebyshev, Gaussian
    - A/B testing
- Kernel density estimation
  - Parzen windows, Nearest neighbors,
     Watson-Nadaraya estimator
- Exponential families
  - Gaussian, multinomial, Poisson
  - Conjugate distributions and smoothing, integrating out



Peninsula Grill

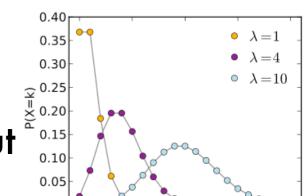


Come check out our new menu specials at your



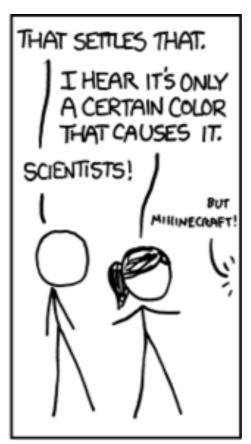
Peninsula Grill

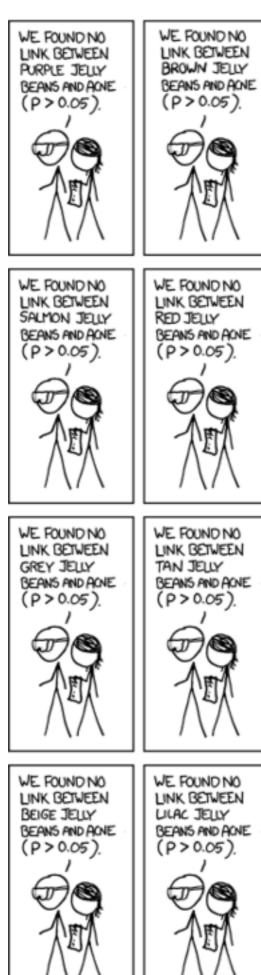
Come check out our new menu specials at your favorite city diner!

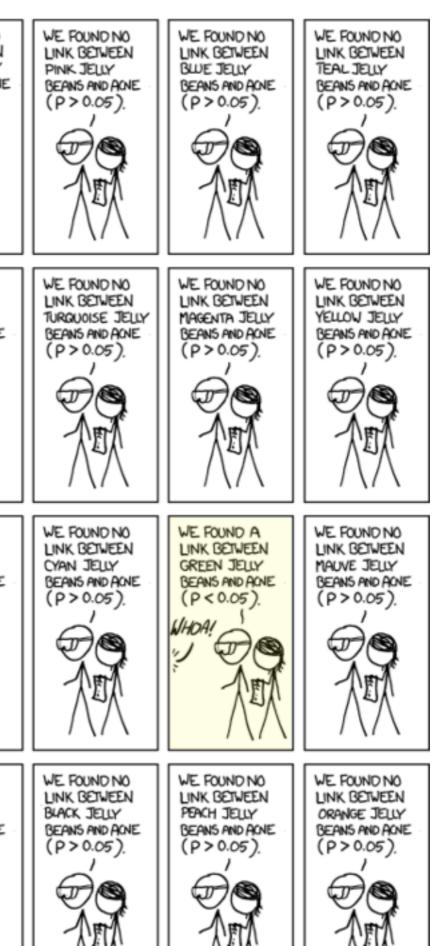


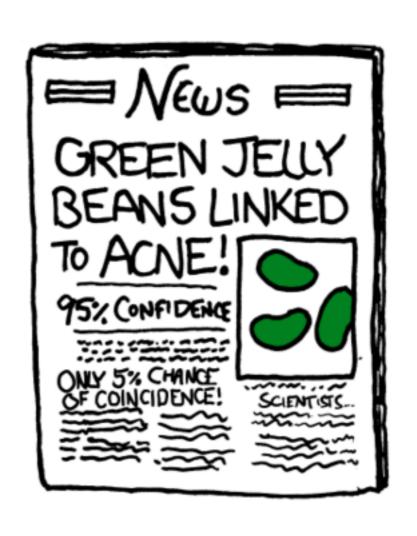












#### 2.1 Probabilities

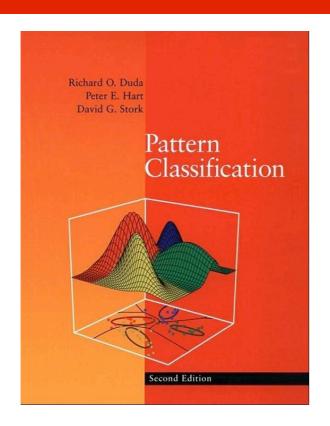


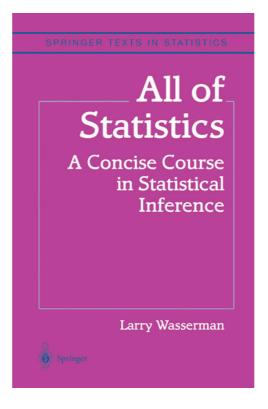
Bayes

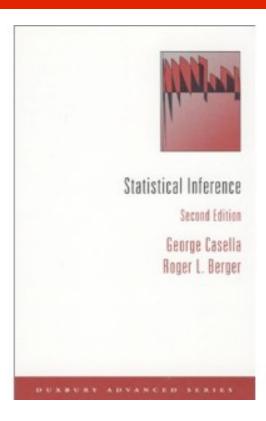


Kolmogorov

#### Statistics 101







# Probability

- Space of events X
  - server working; slow response; server broken
  - income of the user (e.g. \$95,000)
  - query text for search (e.g. "statistics tutorial")
- Probability axioms (Kolmogorov)

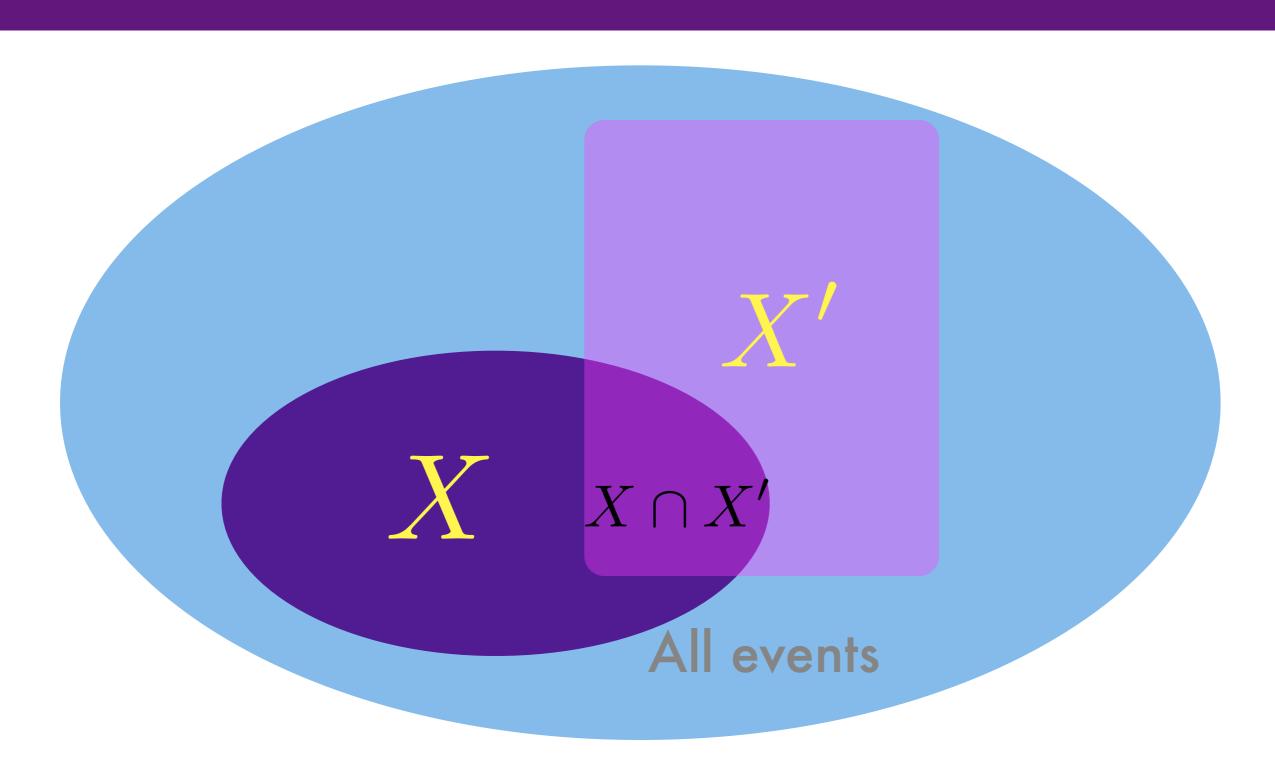
$$\Pr(X) \in [0, 1], \Pr(\mathcal{X}) = 1$$
  
 $\Pr(\bigcup_i X_i) = \sum_i \Pr(X_i) \text{ if } X_i \cap X_j = \emptyset$ 

- Example queries
  - P(server working) = 0.999
  - P(90,000 < income < 100,000) = 0.1

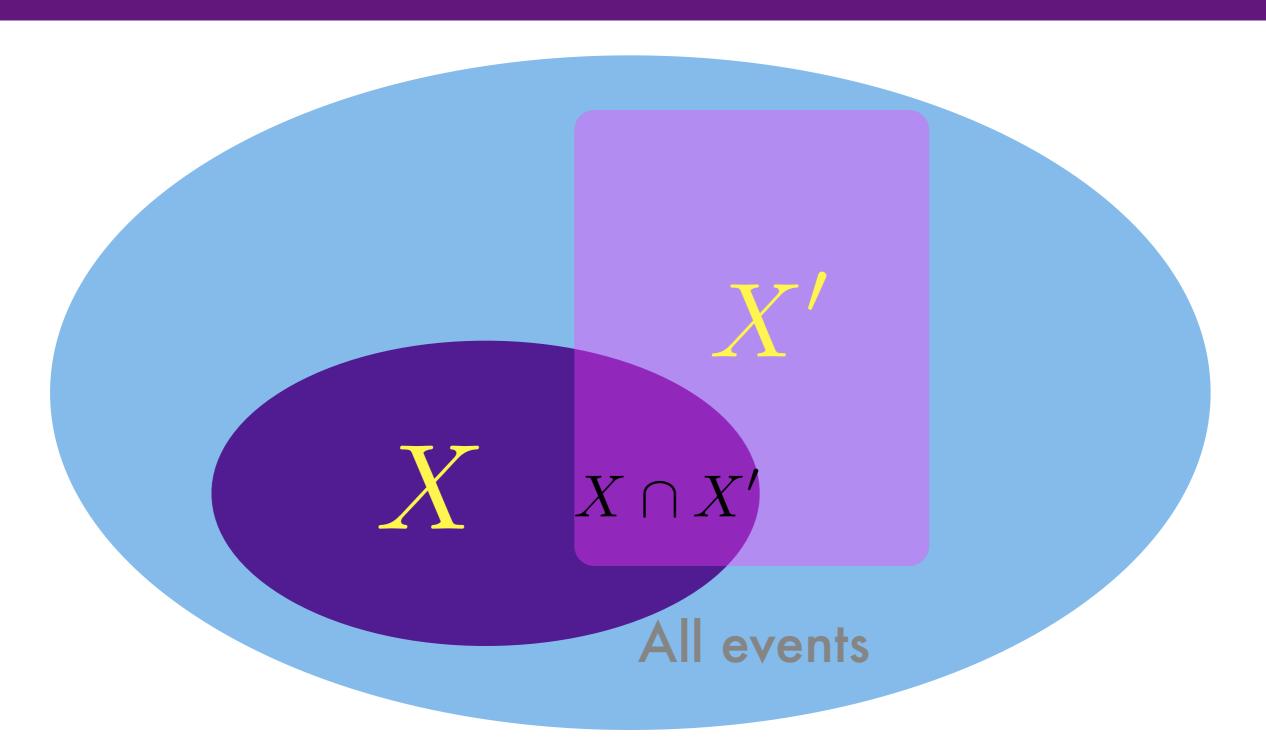
#### Venn Diagram

All events

#### Venn Diagram



#### Venn Diagram



$$Pr(X \cup X') = Pr(X) + Pr(X') - Pr(X \cap X')$$

## (In)dependence

- Independence  $Pr(x, y) = Pr(x) \cdot Pr(y)$ 
  - Login behavior of two users (approximately)
  - Disk crash in different colos (approximately)

# (In)dependence

- Independence  $Pr(x, y) = Pr(x) \cdot Pr(y)$ 
  - Login behavior of two users (approximately)
  - Disk crash in different colos (approximately)
- Dependent events
  - Emails  $\Pr(x,y) \neq \Pr(x) \cdot \Pr(y)$
  - Queries
  - News stream / Buzz / Tweets
  - IM communication
  - Russian Roulette

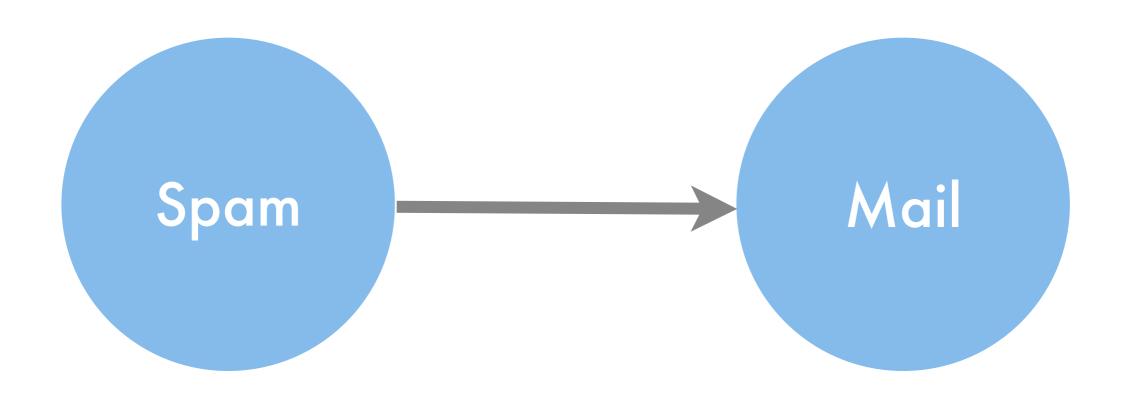
## (In)dependence

- Independence  $Pr(x, y) = Pr(x) \cdot Pr(y)$ 
  - Login behavior of two users (approximately)
  - Disk crash in different colos (approximately)
- Dependent events
  - Emails

$$\Pr(x,y) \neq \Pr(x) \land \Pr(y)$$

- Queries
- News stream / Buzz / Tweets
- IM communication
- Russian Roulette

Everywhere!



p(spam, mail) = p(spam) p(mail|spam)

## Bayes Rule

Joint Probability

$$Pr(X,Y) = Pr(X|Y) Pr(Y) = Pr(Y|X) Pr(X)$$

Bayes Rule

$$Pr(X|Y) = \frac{Pr(Y|X) \cdot Pr(X)}{Pr(Y)}$$

- Hypothesis testing
- Reverse conditioning

# AIDS test (Bayes rule)

- Data
  - Approximately 0.1% are infected
  - Test detects all infections
  - Test reports positive for 1% healthy people
- Probability of having AIDS if test is positive

# AIDS test (Bayes rule)

- Data
  - Approximately 0.1% are infected
  - Test detects all infections
  - Test reports positive for 1% healthy people
- Probability of having AIDS if test is positive

$$\Pr(a = 1|t) = \frac{\Pr(t|a = 1) \cdot \Pr(a = 1)}{\Pr(t)}$$

$$= \frac{\Pr(t|a = 1) \cdot \Pr(a = 1)}{\Pr(t|a = 1) \cdot \Pr(a = 1) + \Pr(t|a = 0) \cdot \Pr(a = 0)}$$

$$= \frac{1 \cdot 0.001}{1 \cdot 0.001 + 0.01 \cdot 0.999} = 0.091$$

- Use a follow-up test
  - Test 2 reports positive for 90% infections
  - Test 2 reports positive for 5% healthy people

$$\frac{0.01 \cdot 0.05 \cdot 0.999}{1 \cdot 0.9 \cdot 0.001 + 0.01 \cdot 0.05 \cdot 0.999} = 0.357$$

- Use a follow-up test
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$$\frac{0.01 \cdot 0.05 \cdot 0.999}{1 \cdot 0.9 \cdot 0.001 + 0.01 \cdot 0.05 \cdot 0.999} = 0.357$$

Why can't we use Test 1 twice?
 Outcomes are not independent but tests 1 and 2 are conditionally independent

- Use a follow-up test
  - Test 2 reports positive for 90% infections
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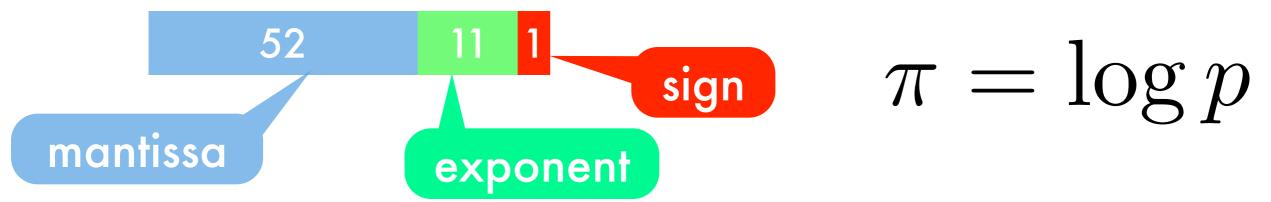
$$\frac{0.01 \cdot 0.05 \cdot 0.999}{1 \cdot 0.9 \cdot 0.001 + 0.01 \cdot 0.05 \cdot 0.999} = 0.357$$

Why can't we use Test 1 twice?
 Outcomes are not independent but tests 1 and 2 are conditionally independent

$$p(t_1, t_2|a) = p(t_1|a) \cdot p(t_2|a)$$

# Logarithms are good

Floating point numbers



- Probabilities can be very small. In particular products of many probabilities. Underflow!
- Store data in mantissa, not exponent

$$\prod_{i} p_{i} \to \sum_{i} \pi_{i} \qquad \sum_{i} p_{i} \to \max \pi + \log \sum_{i} \exp \left[\pi_{i} - \max \pi\right]$$

Known bug e.g. in Mahout Dirichlet clustering

# Application: Naive Bayes



Key assumption
 Words occur independently of each other given the label of the document

$$p(w_1, \dots, w_n | \text{spam}) = \prod_{i=1}^n p(w_i | \text{spam})$$

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 Words occur independently of each other
 given the label of the document

$$p(w_1, \dots, w_n | \text{spam}) = \prod p(w_i | \text{spam})$$

ullet Spam classification via Bayes Rule

$$p(\operatorname{spam}|w_1,\ldots,w_n) \propto p(\operatorname{spam}) \prod_{i=1}^n p(w_i|\operatorname{spam})$$

Key assumption
 Words occur independently of each other
 given the label of the document

$$p(w_1, \dots, w_n | \text{spam}) = \prod p(w_i | \text{spam})$$

• Spam classification via Bayes n

$$p(\operatorname{spam}|w_1,\ldots,w_n) \propto p(\operatorname{spam}) \prod_{i=1}^n p(w_i|\operatorname{spam})$$

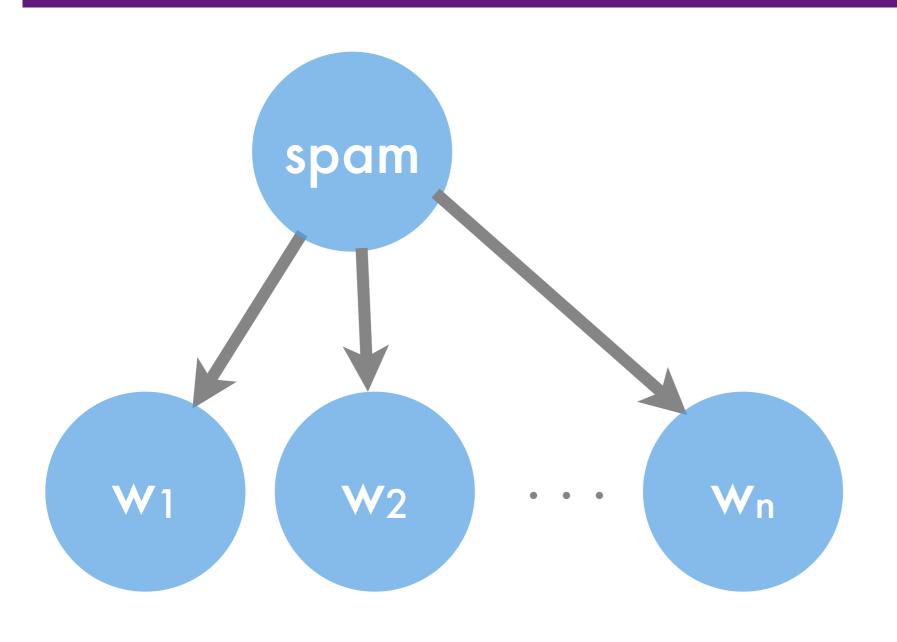
• Parameter estimation Compute spam probability and word distributions for spam and ham

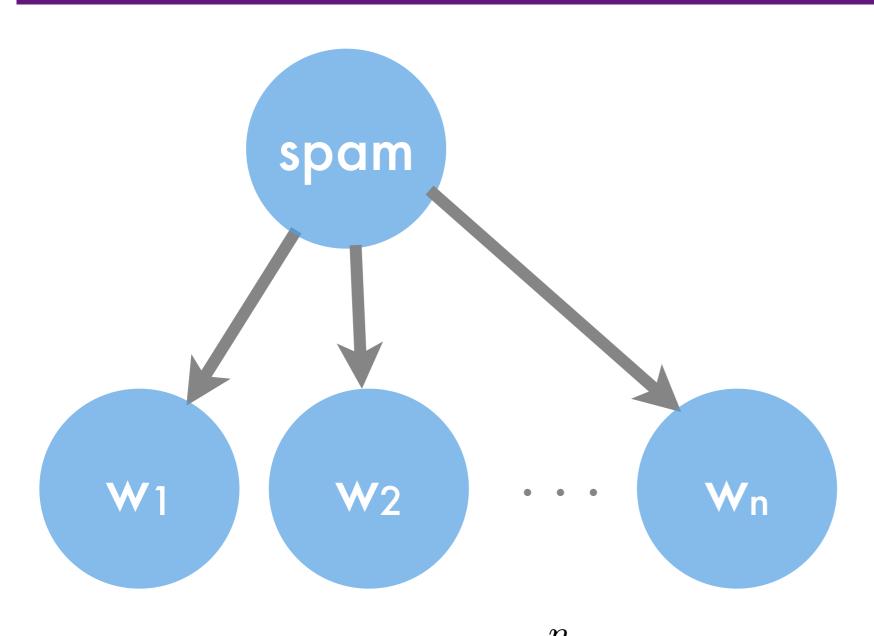
#### Equally likely phrases

- Get rich quick. Buy UCB stock.
- Buy Viagra. Make your UCB experience last longer.
- You deserve a PhD from UCB.
   We recognize your expertise.

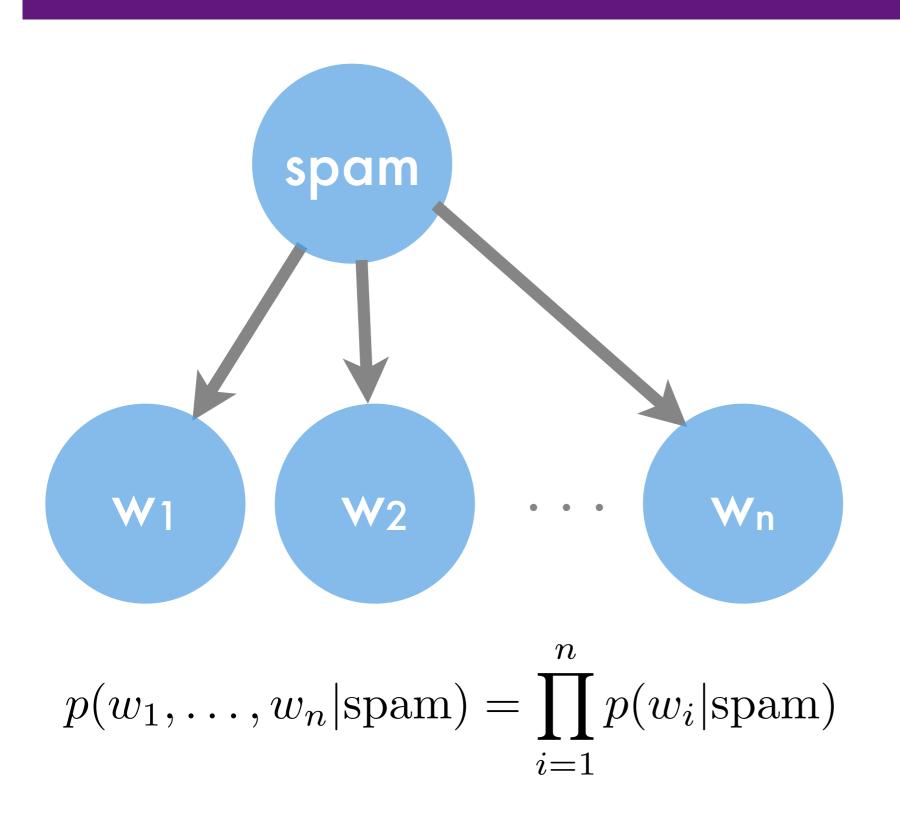
#### Equally likely phrases

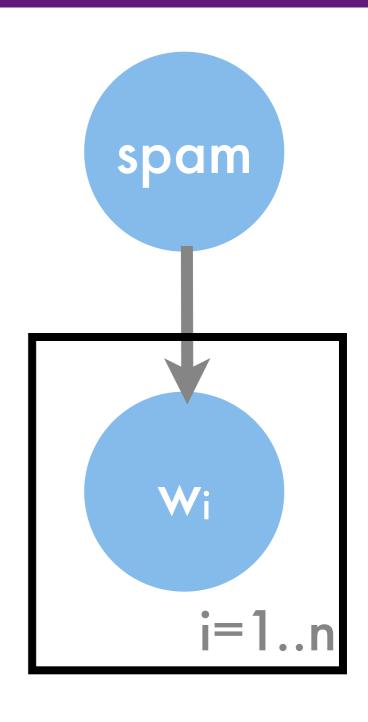
- Get rich quick. Buy UCB stock.
- Buy Viagra. Make your UCB experience last longer.
- You deserve a PhD from UCB.
   We recognize your expertise.
- Make your rich UCB PhD experience last longer.

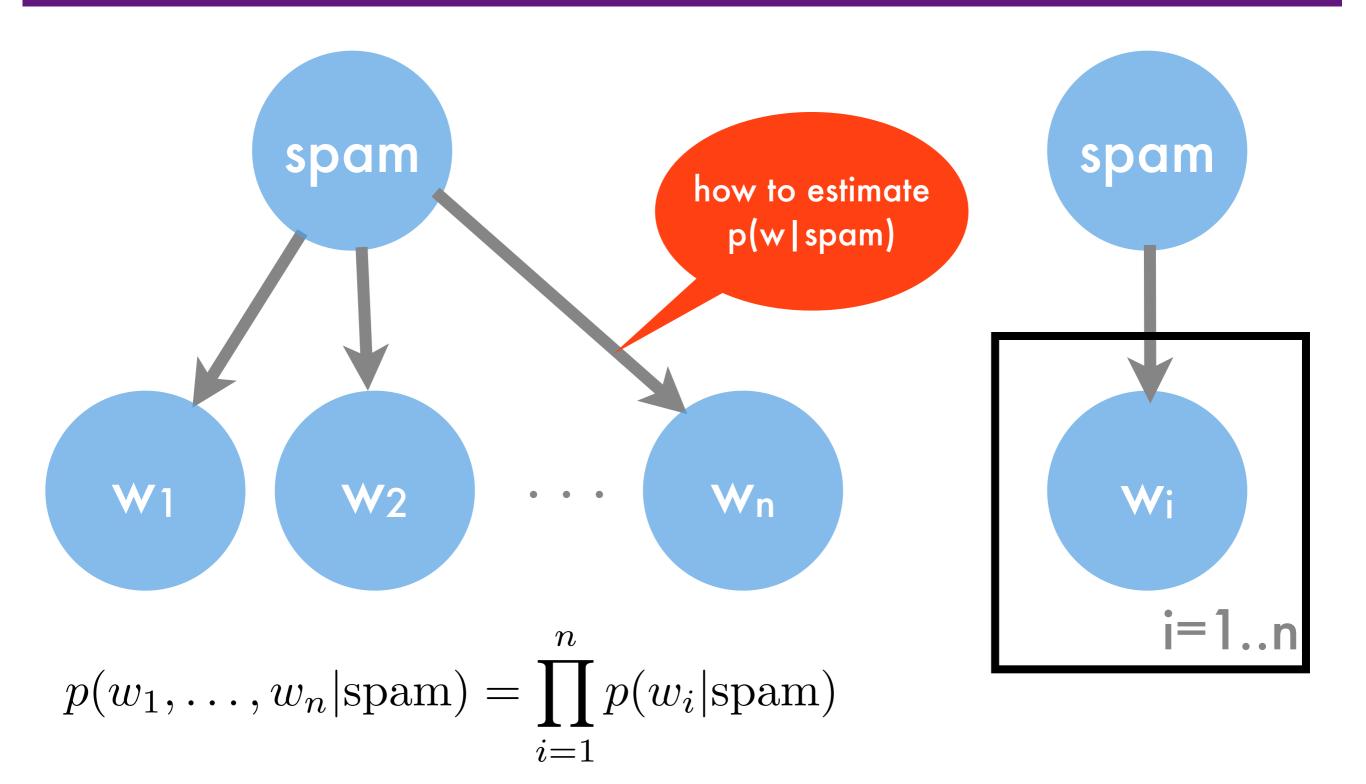




$$p(w_1, \dots, w_n | \text{spam}) = \prod_{i=1}^n p(w_i | \text{spam})$$







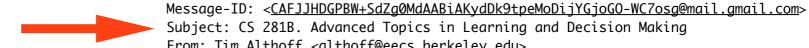
- Data
  - Emails (headers, body, metadata)
  - Labels (spam/ham)
     assume that users actually label all mails
- Processing capability
  - Billions of e-mails
  - 1000s of servers
- Need to estimate p(y), p(xi|y)
  - Compute distribution of x<sub>i</sub> for every y
  - Compute distribution of y

# this is a gross simplification

- date
- time
- recipient path
- IP number
- sender
- encoding
- many more features

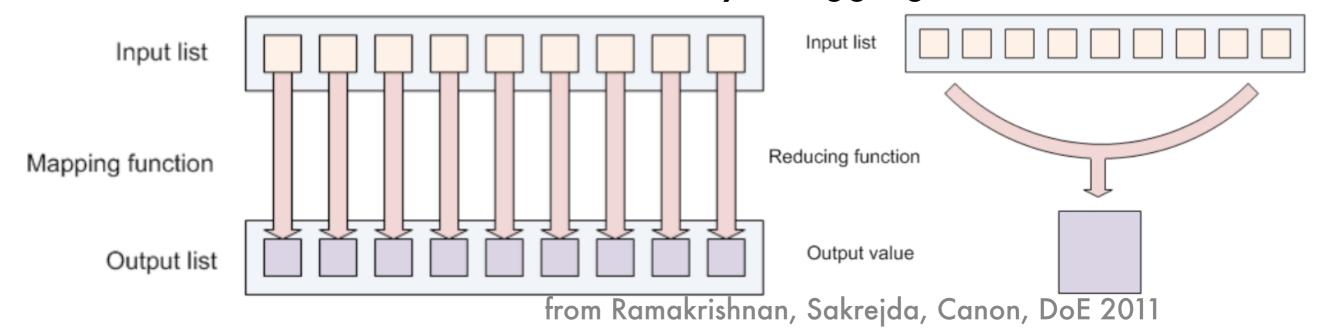
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Received: by 10.213.17.145 with SMTP id s17mr2519891eba.147.1325629071725;
        Tue, 03 Jan 2012 14:17:51 -0800 (PST)
Return-Path: <alex+caf_=alex.smola=<u>gmail.com@smola.ora</u>>
Received: from mail-ey0-f175.google.com (mail-ey0-f175.google.com [209.85.215.175])
        by mx.google.com with ESMTPS id n4si29264232eef.57.2012.01.03.14.17.51
        (version=TLSv1/SSLv3 cipher=OTHER);
        Tue, 03 Jan 2012 14:17:51 -0800 (PST)
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guess record for domain of alex+caf_=alex.smola=<u>gmail.com@smola.org</u>) client-
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permitted nor denied by best guess record for domain of alex
+caf_=alex.smola=<u>gmail.com@smola.ora</u>) smtp.mail=alex+caf_=alex.smola=<u>gmail.com@smola.ora</u>;
dkim=pass (test mode) header.i=@googlemail.com
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        Tue, 03 Jan 2012 14:17:51 -0800 (PST)
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        for <alex@smola.orq>; Tue, 03 Jan 2012 14:17:48 -0800 (PST)
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```

Delivered-To: <a href="mailto:alex.smola@gmail.com">alex.smola@gmail.com</a>



### Recall - Map Reduce

- 1000s of (faulty) machines
- Lots of jobs are mostly embarrassingly parallel (except for a sorting/transpose phase)
- Functional programming origins
  - Map(key,value)
     processes each (key,value) pair and outputs a new (key,value) pair
  - Reduce(key,value)
     reduces all instances with same key to aggregate



## Recall - Map Reduce

- 1000s of (faulty) machines
- Lots of jobs are mostly embarrassingly parallel (except for a sorting/transpose phase)
- Functional programming origins
  - Map(key,value)
     processes each (key,value) pair and outputs a new (key,value) pair
  - Reduce(key,value)
     reduces all instances with same key to aggregate
- Example extremely naive wordcount
  - Map(docID, document)
    for each document emit many (wordID, count) pairs
  - Reduce(wordID, count)
     sum over all counts for given wordID and emit (wordID, aggregate)

## Naive NaiveBayes Classifier

- Two classes (spam/ham)
- Binary features (e.g. presence of \$\$\$, viagra)
- Simplistic Algorithm
  - Count occurrences of feature for spam/ham
  - Count number of spam/ham mails

feature probability

spam probability n(y)

$$p(x_i = \text{TRUE}|y) = \frac{n(i,y)}{n(y)} \text{ and } p(y) = \frac{n(y)}{n}$$

$$p(y|x) \propto \frac{n(y)}{n} \prod_{i:x_i = \text{TRUE}} \frac{n(i,y)}{n(y)} \prod_{i:x_i = \text{FALSE}} \frac{n(y) - n(i,y)}{n(y)}$$

## Naive NaiveBayes Classifier

what if n(i,y)=n(y)?

what if n(i,y)=0?

$$p(y|x) \propto \frac{n(y)}{n} \prod_{i:x_i = \text{TRUE}} \frac{n(i,y)}{n(y)} \prod_{i:x_i = \text{FALSE}} \frac{n(y) - n(i,y)}{n(y)}$$

# Naive NaiveBayes Classifier



what if n(i,y)=0?

$$p(y|x) \propto \frac{n(y)}{n} \prod_{i:x_i = \text{TRUE}} \frac{n(i,y)}{n(y)} \prod_{i:x_i = \text{FALSE}} \frac{n(y) - n(i,y)}{n(y)}$$

# Simple Algorithm

- For each document (x,y) do
  - Aggregate label counts given y
  - For each feature x<sub>i</sub> in x do
    - Aggregate statistic for (x<sub>i</sub>, y) for each y
- For y estimate distribution p(y)
- For each (x<sub>i</sub>,y) pair do
   Estimate distribution p(x<sub>i</sub>|y), e.g. Parzen Windows,
   Exponential family (Gauss, Laplace, Poisson, ...), Mixture
- Given new instance compute

$$p(y|x) \propto p(y) \prod_{i} p(x_{i}|y)$$

# Simple Algorithm

- For each document (x,y) do
  - Aggregate label counts given y pass over all data
  - For each feature x<sub>i</sub> in x do
    - Aggregate statistic for (x<sub>i</sub>, y) for each y
- For y estimate distribution p(y)
- For each (x<sub>i</sub>,y) pair do
   Estimate distribution p(x<sub>i</sub>|y), e.g. Parzen Windows,
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- Given new instance compute

$$p(y|x) \propto p(y) \prod_{i} p(x_{i}|y)$$

## MapReduce Algorithm

- Map(document (x,y))
  - For each mapper for each feature  $x_i$  in x do
    - Aggregate statistic for (xi, y) for each y
  - Send statistics (key =  $(x_i,y)$ , value = counts) to reducer
- Reduce(x<sub>i</sub>, y)
  - Aggregate over all messages from mappers
  - Estimate distribution  $p(x_i|y)$ , e.g. Parzen Windows, Exponential family (Gauss, Laplace, Poisson, ...), Mixture
  - Send coordinate-wise model to global storage
- Given new instance compute

$$p(y|x) \propto p(y) \prod_{i} p(x_{i}|y)$$

## MapReduce Algorithm

- Map(document (x,y))
  - For each mapper for each feature x<sub>i</sub> in x do local per
    - Aggregate statistic for (xi, y) for each y chunkserver
  - Send statistics (key =  $(x_i,y)$ , value = counts) to reducer
- Reduce(x<sub>i</sub>, y)

- only aggregates
- Aggregate over all messages from mappers needed
- Estimate distribution p(x<sub>i</sub>|y), e.g. Parzen Windows,
   Exponential family (Gauss, Laplace, Poisson, ...), Mixture
- Send coordinate-wise model to global storage
- Given new instance compute

$$p(y|x) \propto p(y) \prod_{i} p(x_{i}|y)$$

### Estimating Probabilities



### Binomial Distribution

- Two outcomes (head, tail); (0,1)
- Data likelihood

$$p(X;\pi) = \pi^{n_1} (1-\pi)^{n_0}$$





- Incorporate constraint via  $p(x;\theta) = \frac{e^{x\theta}}{1 + e^{\theta}}$
- Taking derivatives yields

$$\theta = \log \frac{n_1}{n_0 + n_1} \iff p(x = 1) = \frac{n_1}{n_0 + n_1}$$



### ... in detail ...

$$p(X;\theta) = \prod_{i=1}^{n} p(x_i;\theta) = \prod_{i=1}^{n} \frac{e^{\theta x_i}}{1 + e^{\theta}}$$

$$\Longrightarrow \log p(X;\theta) = \theta \sum_{i=1}^{n} x_i - n \log \left[1 + e^{\theta}\right]$$

$$\Longrightarrow \partial_{\theta} \log p(X;\theta) = \sum_{i=1}^{n} x_i - n \frac{e^{\theta}}{1 + e^{\theta}}$$

$$\Longleftrightarrow \frac{1}{n} \sum_{i=1}^{n} x_i = \frac{e^{\theta}}{1 + e^{\theta}} = p(x = 1)$$

### ... in detail ...

$$p(X;\theta) = \prod_{i=1}^{n} p(x_i;\theta) = \prod_{i=1}^{n} \frac{e^{\theta x_i}}{1 + e^{\theta}}$$

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$$\Longrightarrow \partial_{\theta} \log p(X;\theta) = \sum_{i=1}^{n} x_i - n \frac{e^{\theta}}{1 + e^{\theta}}$$

$$\Longleftrightarrow \frac{1}{n} \sum_{i=1}^{n} x_i = \frac{e^{\theta}}{1 + e^{\theta}} = p(x = 1)$$

empirical probability of x=1

### Discrete Distribution

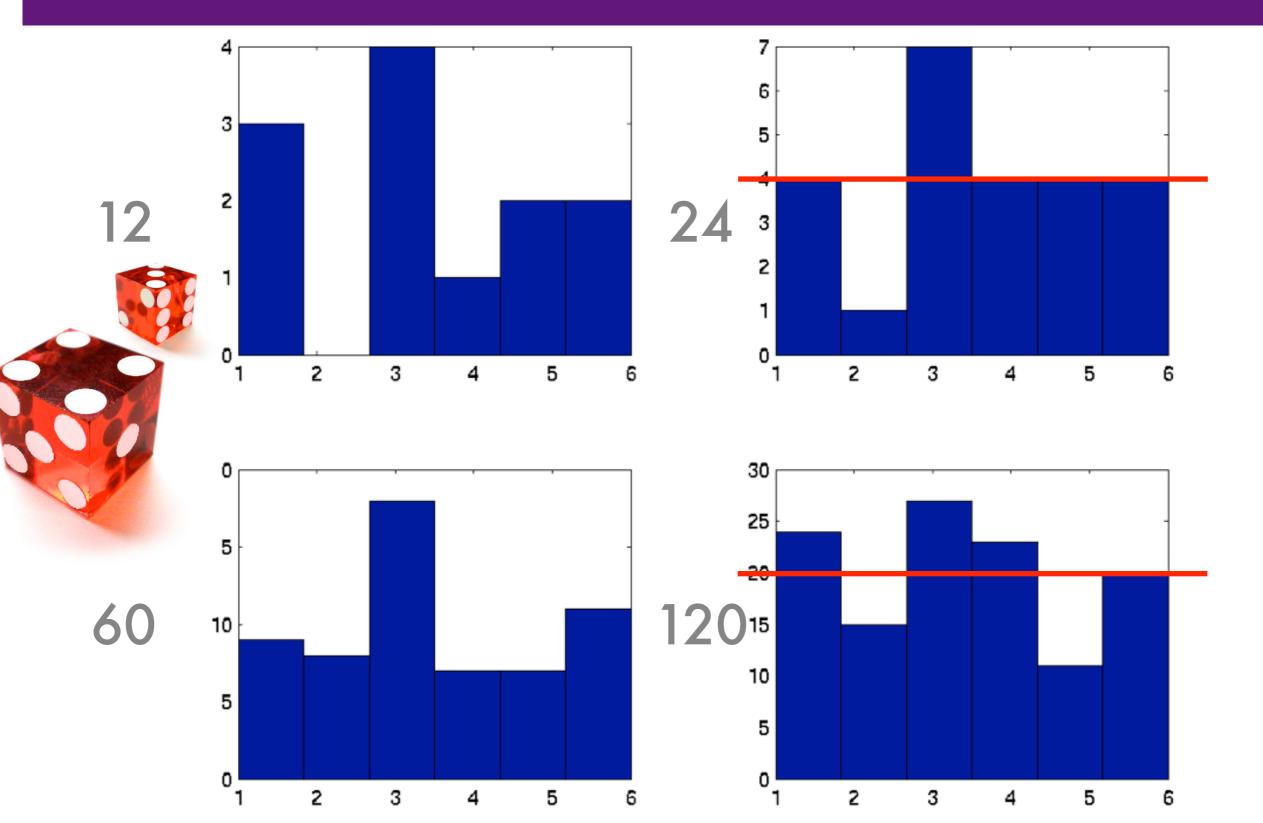
- n outcomes (e.g. USA, Canada, India, UK, NZ)
- Data likelihood

$$p(X;\pi) = \prod \pi_i^{n_i}$$

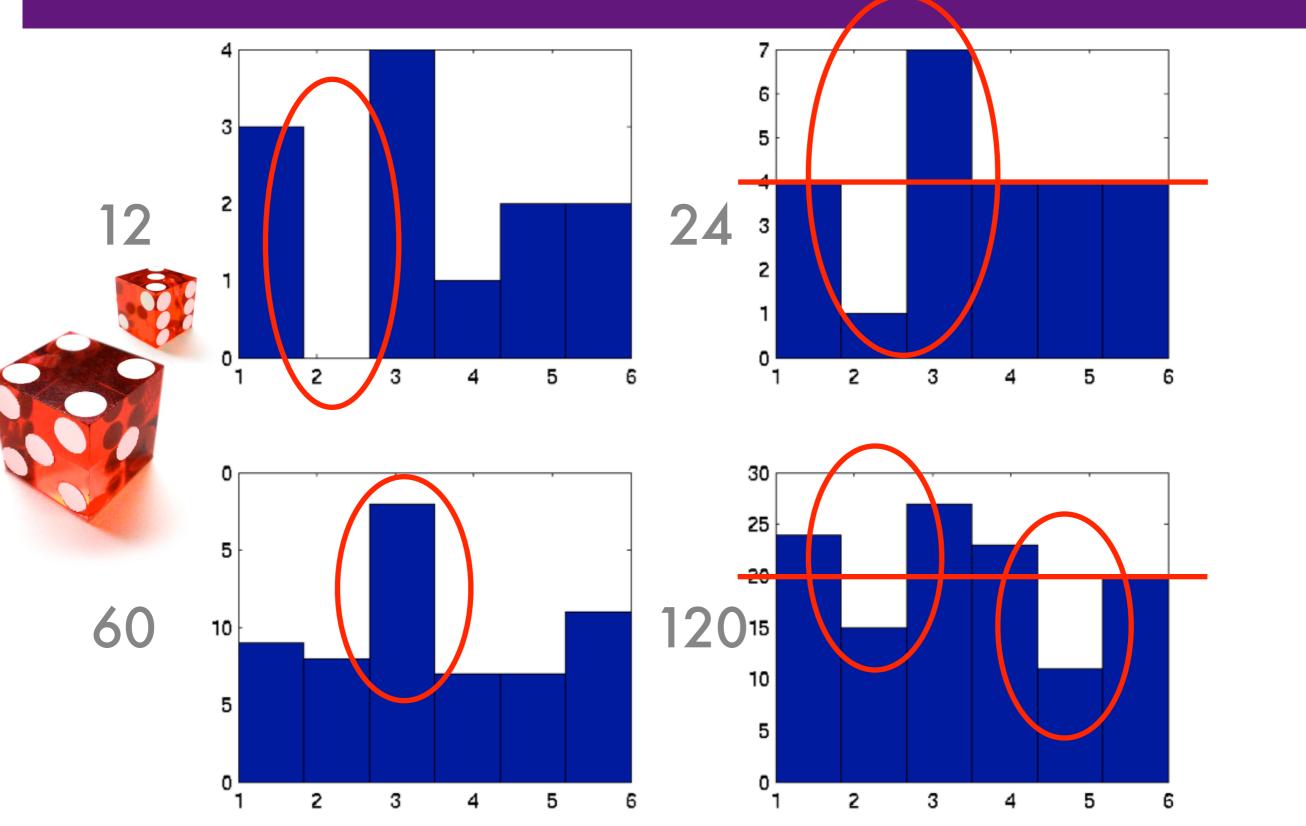
- Maximum Likelihood Estimation
  - Constrained optimization problem ... or ...
  - Incorporate constraint via  $p(x;\theta) = \frac{\exp \theta_x}{\sum_{x'} \exp \theta_{x'}}$
  - Taking derivatives yields

$$\theta_i = \log \frac{n_i}{\sum_j n_j} \iff p(x = i) = \frac{n_i}{\sum_j n_j}$$

# Tossing a Dice



# Tossing a Dice



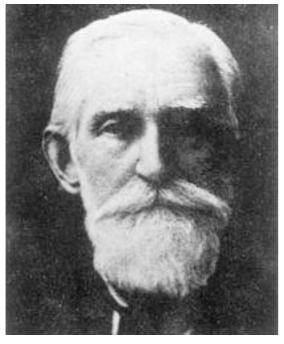
### Key Questions

- Do empirical averages converge?
  - Probabilities
  - Means / moments
- Rate of convergence and limit distribution
- Worst case guarantees
- Using prior knowledge

drug testing, semiconductor fabs computational advertising user interface design ...



### 2.2 Tail Bounds



Chebyshev



Chernoff



Hoeffding

### Expectations

- Random variable x with probability measure
- Expected value of f(x)

$$\mathbf{E}[f(x)] = \int f(x)dp(x)$$

• Special case - discrete probability mass

$$\Pr\{x = c\} = \mathbf{E}[\{x = c\}] = \int \{x = c\} \, dp(x)$$

(same trick works for intervals)

- Draw x<sub>i</sub> identically and independently from p
- Empirical average

$$\mathbf{E}_{\text{emp}}[f(x)] = \frac{1}{n} \sum_{i=1}^{n} f(x_i) \text{ and } \Pr_{\text{emp}} \{x = c\} = \frac{1}{n} \sum_{i=1}^{n} \{x_i = c\}$$

### Deviations

Gambler rolls dice 100 times

$$\hat{P}(X=6) = \frac{1}{n} \sum_{i=1}^{n} \{x_i = 6\}$$

• '6' only occurs 11 times. Fair number is 16.7

#### IS THE DICE TAINTED?

Probability of seeing '6' at most 11 times

$$\Pr(X \le 11) = \sum_{i=0}^{11} p(i) = \sum_{i=0}^{11} {100 \choose i} \left[\frac{1}{6}\right]^i \left[\frac{5}{6}\right]^{100-i} \approx 7.0\%$$

It's probably OK ... can we develop general theory?

#### Deviations

Gambler rolls dice 100 times

$$\hat{P}(X=6) = \frac{1}{n} \sum_{i=1}^{n} \{x_i = 6\}$$

• '6' only occurs 11 times. Fair number is 16.7

IS THE DICE TAINTED?

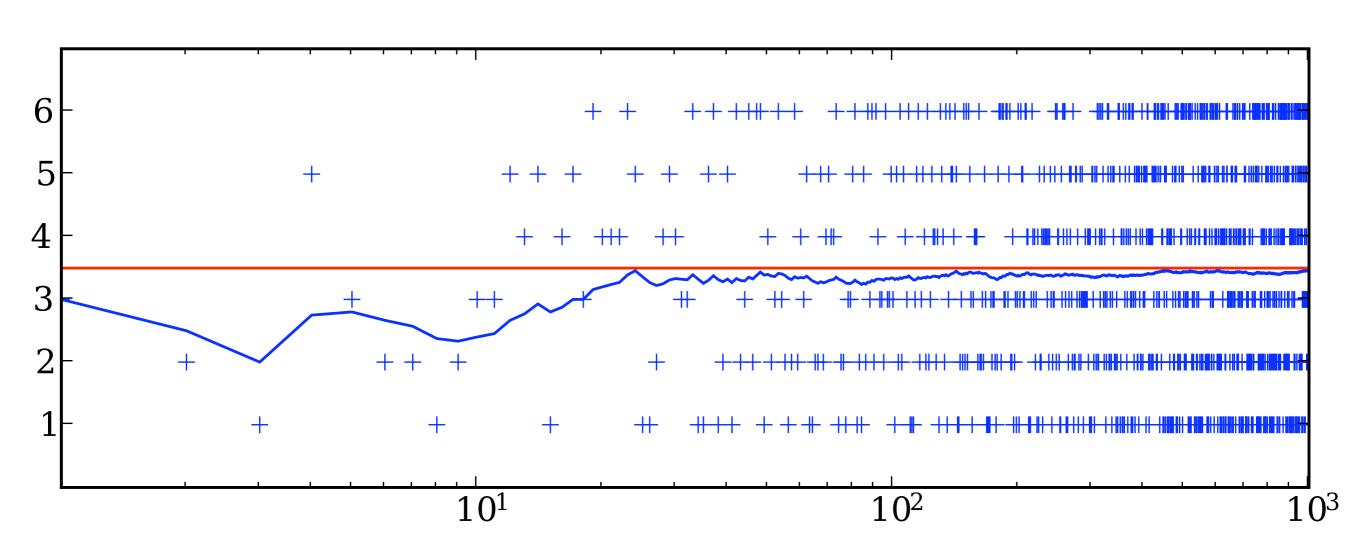
ad campaign working new page layout better drug working

Probability of seeing '6' at most 11 times

$$\Pr(X \le 11) = \sum_{i=0}^{11} p(i) = \sum_{i=0}^{11} {100 \choose i} \left[\frac{1}{6}\right]^i \left[\frac{5}{6}\right]^{100-i} \approx 7.0\%$$

It's probably OK ... can we develop general theory?

### Empirical average for a dice



how quickly does it converge?

# Law of Large Numbers

- Random variables  $x_i$  with mean  $\mu = \mathbf{E}[x_i]$
- Empirical average  $\hat{\mu}_n := \frac{1}{n} \sum_{i=1}^n x_i$
- Weak Law of Large Numbers

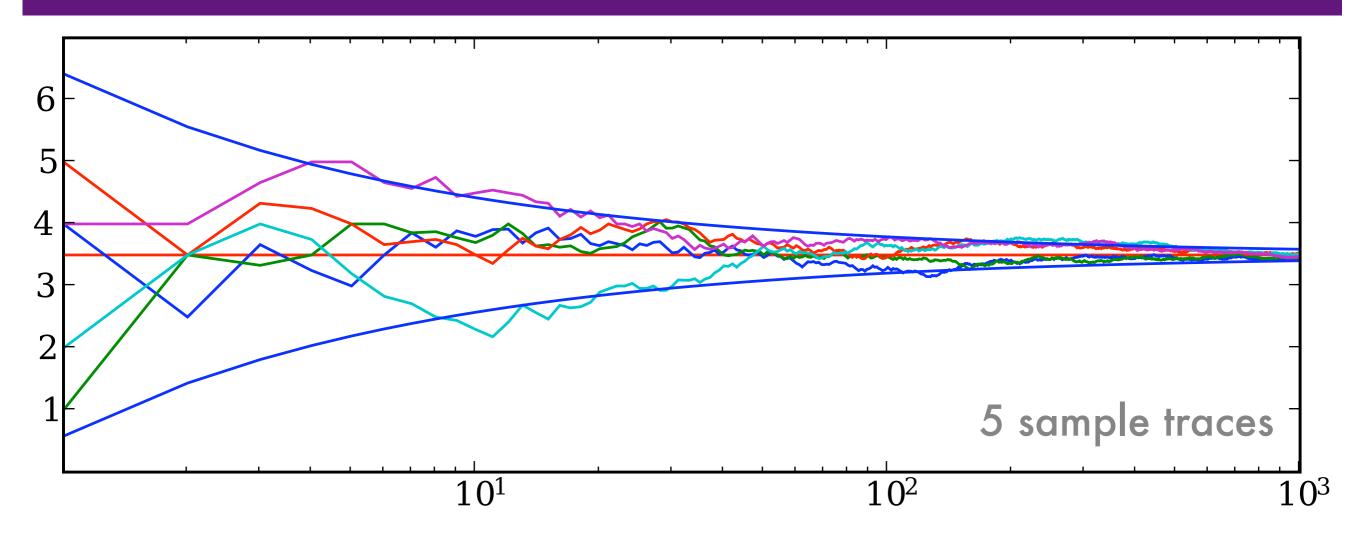
$$\lim_{n \to \infty} \Pr(|\hat{\mu}_n - \mu| \le \epsilon) = 1 \text{ for any } \epsilon > 0$$

Strong Law of Large Numbers

$$\Pr\left(\lim_{n\to\infty}\hat{\mu}_n = \mu\right) = 1$$

this means convergence in probability

### Empirical average for a dice



- Upper and lower bounds are  $\mu \pm \sqrt{{\rm Var}(x)/n}$
- This is an example of the central limit theorem

### Central Limit Theorem

- Independent random variables  $x_i$  with mean  $\mu_i$  and standard deviation  $\sigma_i$
- The random variable

$$z_n := \left[\sum_{i=1}^n \sigma_i^2\right]^{-\frac{1}{2}} \left[\sum_{i=1}^n x_i - \mu_i\right]$$

converges to a Normal Distribution  $\mathcal{N}(0,1)$ 

### Central Limit Theorem

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converges to a Normal Distribution  $\mathcal{N}(0,1)$ 

Special case - IID random variables & average

$$\frac{\sqrt{n}}{\sigma} \left[ \frac{1}{n} \sum_{i=1}^{n} x_i - \mu \right] \to \mathcal{N}(0,1)$$

$$O\left(n^{-\frac{1}{2}}\right)$$
 convergence

### Slutsky's Theorem

- Continuous mapping theorem
  - X<sub>i</sub> and Y<sub>i</sub> sequences of random variables
  - X<sub>i</sub> has as its limit the random variable X
  - Yi has as its limit the constant c
  - g(x,y) is continuous function for all g(x,c)

g(X<sub>i</sub>, Y<sub>i</sub>) converges in distribution to g(X,c)

### Delta Method

Random variable X<sub>i</sub> convergent to b

$$a_n^{-2}(X_n - b) \to \mathcal{N}(0, \Sigma)$$
 with  $a_n^2 \to 0$  for  $n \to \infty$ 

- g is a continuously differentiable function for b
- Then g(Xi) inherits convergence properties

$$a_n^{-2}\left(g(X_n) - g(b)\right) \to \mathcal{N}(0, \left[\nabla_x g(b)\right] \Sigma \left[\nabla_x g(b)\right]^\top\right)$$

Proof: use Taylor expansion for g(X<sub>n</sub>) - g(b)

$$a_n^{-2} [g(X_n) - g(b)] = [\nabla_x g(\xi_n)]^{\top} a_n^{-2} (X_n - b)$$

- $g(\xi_n)$  is on line segment  $[X_n, b]$
- By Slutsky's theorem it converges to g(b)
- Hence g(X<sub>i</sub>) is asymptotically normal



# Tools for the proof

### Fourier Transform

Fourier transform relations

$$F[f](\omega) := (2\pi)^{-\frac{d}{2}} \int_{\mathbb{R}^n} f(x) \exp(-i \langle \omega, x \rangle) dx$$
$$F^{-1}[g](x) := (2\pi)^{-\frac{d}{2}} \int_{\mathbb{R}^n} g(\omega) \exp(i \langle \omega, x \rangle) d\omega.$$

- Useful identities
  - Identity

$$F^{-1} \circ F = F \circ F^{-1} = \mathrm{Id}$$

Derivative

$$F[\partial_x f] = -i\omega F[f]$$

Convolution (also holds for inverse transform)

$$F[f \circ g] = (2\pi)^{\frac{d}{2}} F[f] \cdot F[g]$$

#### The Characteristic Function Method

Characteristic function

$$\phi_X(\omega) := F^{-1}[p(x)] = \int \exp(i\langle \omega, x \rangle) dp(x)$$

- For X and Y independent we have
  - Joint distribution is convolution

$$p_{X+Y}(z) = \int p_X(z-y)p_Y(y)dy = p_X \circ p_Y$$

Characteristic function is product

$$\phi_{X+Y}(\omega) = \phi_X(\omega) \cdot \phi_Y(\omega)$$

- Proof plug in definition of Fourier transform
- Characteristic function is unique

### Proof - Weak law of large numbers

- Require that expectation exists
- Taylor expansion of exponential

$$\exp(iwx) = 1 + i\langle w, x \rangle + o(|w|)$$
  
and hence  $\phi_X(\omega) = 1 + iw\mathbf{E}_X[x] + o(|w|).$ 

(need to assume that we can bound the tail)

Average of random variables

$$\phi_{\hat{\mu}_m}(\omega) = \left(1 + \frac{i}{m}w\mu + o(m^{-1}|w|)\right)^m$$

convolution

Limit is constant distribution

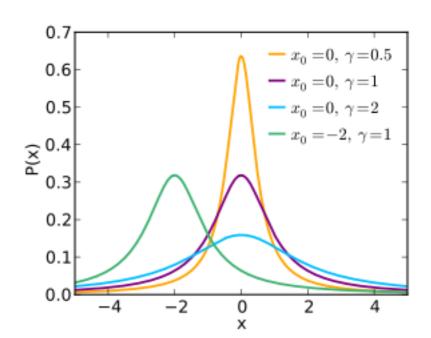
$$\phi_{\hat{\mu}_m}(\omega) \to \exp i\omega\mu = 1 + i\omega\mu + \dots$$

vanishing higher order terms

## Warning

- Moments may not always exist
  - Cauchy distribution

$$p(x) = \frac{1}{\pi} \frac{1}{1 + x^2}$$



 For the mean to exist the following integral would have to converge

$$\int |x| dp(x) \ge \frac{2}{\pi} \int_{1}^{\infty} \frac{x}{1+x^2} dx \ge \frac{1}{\pi} \int_{1}^{\infty} \frac{1}{x} dx = \infty$$

### Proof - Central limit theorem

- Require that second order moment exists (we assume they're all identical WLOG)
- Characteristic function

$$\exp(iwx) = 1 + iwx - \frac{1}{2}w^2x^2 + o(|w|^2)$$

and hence 
$$\phi_X(\omega) = 1 + iw \mathbf{E}_X[x] - \frac{1}{2} w^2 \text{var}_X[x] + o(|w|^2)$$

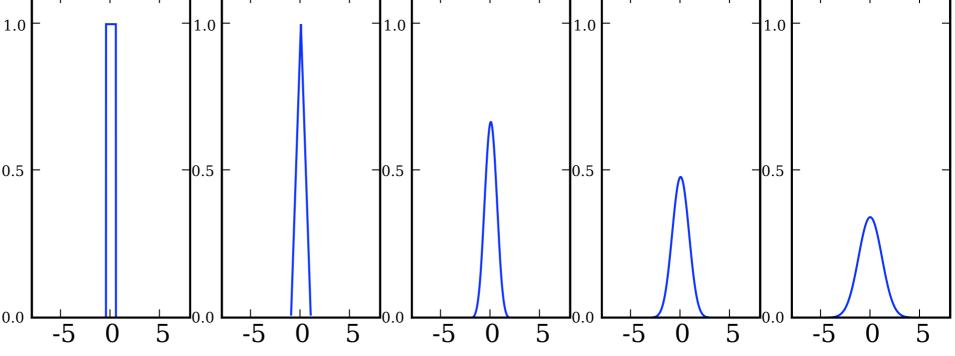
• Subtract out mean (centering)  $z_n \coloneqq \left[\sum_{i=1}^n \sigma_i^2\right]^{-\frac{1}{2}} \left[\sum_{i=1}^n x_i - \mu_i\right]$ 

$$\phi_{Z_m}(\omega) = \left(1 - \frac{1}{2m}w^2 + o(m^{-1}|w|^2)\right)^m \to \exp\left(-\frac{1}{2}w^2\right) \text{ for } m \to \infty$$

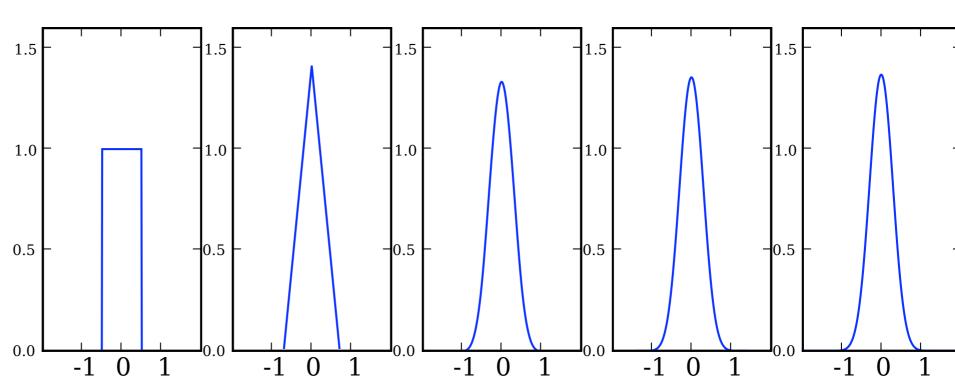
This is the FT of a Normal Distribution

#### Central Limit Theorem in Practice





scaled



# Finite sample tail bounds



# Simple tail bounds

Gauss Markov inequality

Random variable X with mean  $\mu$ 

$$\Pr(X \ge \epsilon) \le \mu/\epsilon$$

Proof - decompose expectation

$$\Pr(X \ge \epsilon) = \int_{\epsilon}^{\infty} dp(x) \le \int_{\epsilon}^{\infty} \frac{x}{\epsilon} dp(x) \le \epsilon^{-1} \int_{0}^{\infty} x dp(x) = \frac{\mu}{\epsilon}.$$

Chebyshev inequality

Random variable X with mean  $\mu$  and variance  $\sigma^2$ 

$$\Pr(|\hat{\mu}_m - \mu| > \epsilon) \le \sigma^2 m^{-1} \epsilon^{-2}$$
 or equivalently  $\epsilon \le \sigma / \sqrt{m\delta}$ 

Proof - applying Gauss-Markov to  $Y = (X - \mu)^2$  with confidence  $\epsilon^2$  yields the result.

# Scaling behavior

Gauss-Markov

$$\epsilon \leq \frac{\mu}{\delta}$$

Scales properly in  $\mu$  but expensive in  $\delta$ 

Chebyshev

$$\epsilon \le \frac{\sigma}{\sqrt{m\delta}}$$

Proper scaling in  $\sigma$  but still bad in  $\delta$ 

Can we get logarithmic scaling in  $\delta$ ?

### Chernoff bound

KL-divergence variant of Chernoff bound

$$K(p,q) = p \log \frac{p}{q} + (1-p) \log \frac{1-p}{1-q}$$

• n independent tosses from biased coin with p

$$\Pr\left\{\sum_{i} x_i \ge nq\right\} \le \exp\left(-nK(q,p)\right) \le \exp\left(-2n(p-q)^2\right)$$

Pinsker's inequality

• Proof w.l.o.g.q > p and set  $k \ge qn$ 

$$\frac{\Pr\left\{\sum_{i} x_{i} = k | q\right\}}{\Pr\left\{\sum_{i} x_{i} = k | p\right\}} = \frac{q^{k} (1 - q)^{n - k}}{p^{k} (1 - p)^{n - k}} \ge \frac{q^{qn} (1 - q)^{n - qn}}{p^{qn} (1 - p)^{n - qn}} = \exp\left(nK(q, p)\right)$$

$$\sum_{k \ge nq} \Pr\left\{\sum_{i} x_i = k | p\right\} \le \sum_{k \ge nq} \Pr\left\{\sum_{i} x_i = k | q\right\} \exp(-nK(q, p)) \le \exp(-nK(q, p))$$

# McDiarmid Inequality

- Independent random variables X<sub>i</sub>
- Function  $f: \mathcal{X}^m \to \mathbb{R}$
- Deviation from expected value

$$\Pr\left(|f(x_1,\ldots,x_m)-\mathbf{E}_{X_1,\ldots,X_m}[f(x_1,\ldots,x_m)]|>\epsilon\right)\leq 2\exp\left(-2\epsilon^2C^{-2}\right)$$
Here C is given by  $C^2=\sum_{i=1}^m c_i^2$  where

$$|f(x_1,\ldots,x_i,\ldots,x_m)-f(x_1,\ldots,x_i',\ldots,x_m)| \le c_i$$

Hoeffding's theorem
 f is average and X<sub>i</sub> have bounded range c

$$\Pr(|\hat{\mu}_m - \mu| > \epsilon) \le 2 \exp\left(-\frac{2m\epsilon^2}{c^2}\right).$$

# Scaling behavior

#### Hoeffding

$$\delta := \Pr\left(|\hat{\mu}_m - \mu| > \epsilon\right) \le 2 \exp\left(-\frac{2m\epsilon^2}{c^2}\right)$$

$$\Longrightarrow \log \delta/2 \le -\frac{2m\epsilon^2}{c^2}$$

$$\Longrightarrow \epsilon \le c\sqrt{\frac{\log 2 - \log \delta}{2m}}$$

This helps when we need to combine several tail bounds since we only pay logarithmically in terms of their combination.

#### More tail bounds

- Higher order moments
  - Bernstein inequality (needs variance bound)

$$\Pr\left(\mu_m - \mu \ge \epsilon\right) \le \exp\left(-\frac{t^2/2}{\sum_i \mathbf{E}[X_i^2] + Mt/3}\right)$$

here M upper-bounds the random variables Xi

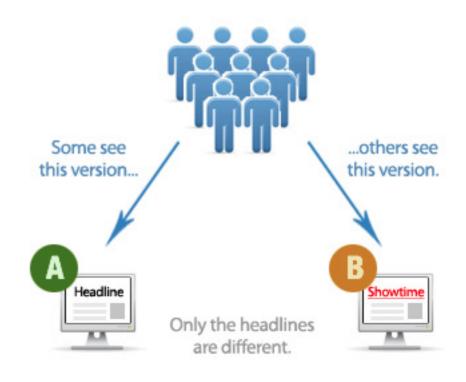
- Proof via Gauss-Markov inequality applied to exponential sums (hence exp. inequality)
- See also Azuma, Bennett, Chernoff, ...
- Absolute / relative error bounds
- Bounds for (weakly) dependent random variables

# Tail bounds in practice



# A/B testing

- Two possible webpage layouts
- Which layout is better?
- Experiment
  - Half of the users see A
  - The other half sees design B



 How many trials do we need to decide which page attracts more clicks?

Assume that the probabilities are p(A) = 0.1 and p(B) = 0.11 respectively and that p(A) is known

# Chebyshev Inequality

- Need to bound for a deviation of 0.01
- Mean is p(B) = 0.11 (we don't know this yet)
- Want failure probability of 5%
- If we have no prior knowledge, we can only bound the variance by  $\sigma^2 = 0.25$

$$m \le \frac{\sigma^2}{\epsilon^2 \delta} = \frac{0.25}{0.01^2 \cdot 0.05} = 50,000$$

• If we know that the click probability is at most 0.15 we can bound the variance at 0.15 \* 0.85 = 0.1275. This requires at most 25,500 users.

# Hoeffding's bound

- Random variable has bounded range [0, 1] (click or no click), hence c=1
- Solve Hoeffding's inequality for m

$$m \le -\frac{c^2 \log \delta/2}{2\epsilon^2} = -\frac{1 \cdot \log 0.025}{2 \cdot 0.01^2} < 18,445$$

This is slightly better than Chebyshev.

# Normal Approximation (Central Limit Theorem)

- Use asymptotic normality
- Gaussian interval containing 0.95 probability

$$\frac{1}{2\pi\sigma^2} \int_{\mu-\epsilon}^{\mu+\epsilon} \exp\left(-\frac{(x-\mu)^2}{2\sigma^2}\right) dx = 0.95$$

is given by  $\epsilon = 2.96\sigma$ .

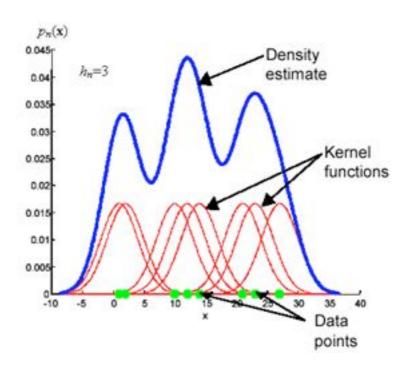
Use variance bound of 0.1275 (see Chebyshev)

$$m \le \frac{2.96^2 \sigma^2}{\epsilon^2} = \frac{2.96^2 \cdot 0.1275}{0.01^2} \le 11,172$$

Same rate as Hoeffding bound! Better bounds by bounding the variance.

# Beyond

- Many different layouts?
- Combinatorial strategy to generate them (aka the Thai Restaurant process)
- What if it depends on the user / time of day
- Stateful user (e.g. query keywords in search)
- What if we have a good prior of the response (rather than variance bound)?
- Explore/exploit/reinforcement learning/control (more details at the end of this class)



### 2.3 Kernel Density Estimation



Parzen

# Density Estimation

- For discrete bins (e.g. male/female; English/French/German/Spanish/Chinese) we get good uniform convergence:
  - Applying the union bound and Hoeffding

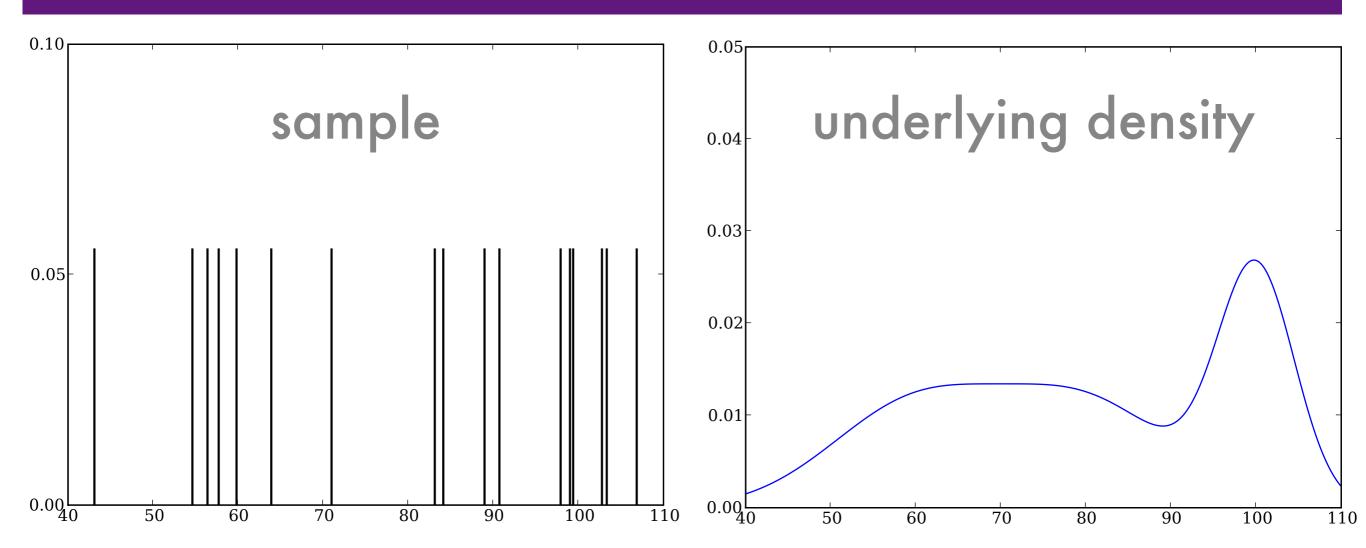
$$\Pr\left(\sup_{a \in A} |\hat{p}(a) - p(a)| \ge \epsilon\right) \le \sum_{a \in A} \Pr\left(|\hat{p}(a) - p(a)| \ge \epsilon\right)$$
$$\le 2|A| \exp\left(-2m\epsilon^2\right)$$

Solving for error probability

good news

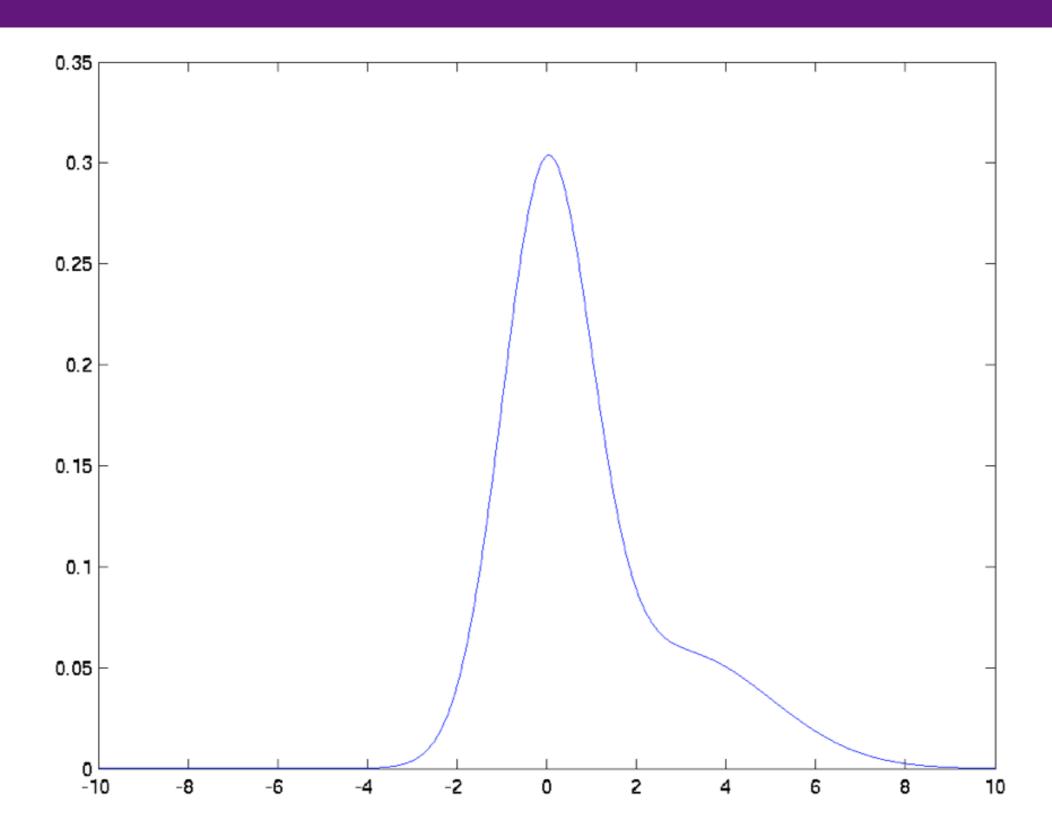
$$\frac{\delta}{2|A|} \le \exp(-m\epsilon^2) \Longrightarrow \epsilon \le \sqrt{\frac{\log 2|A| - \log \delta}{2m}}$$

## Density Estimation

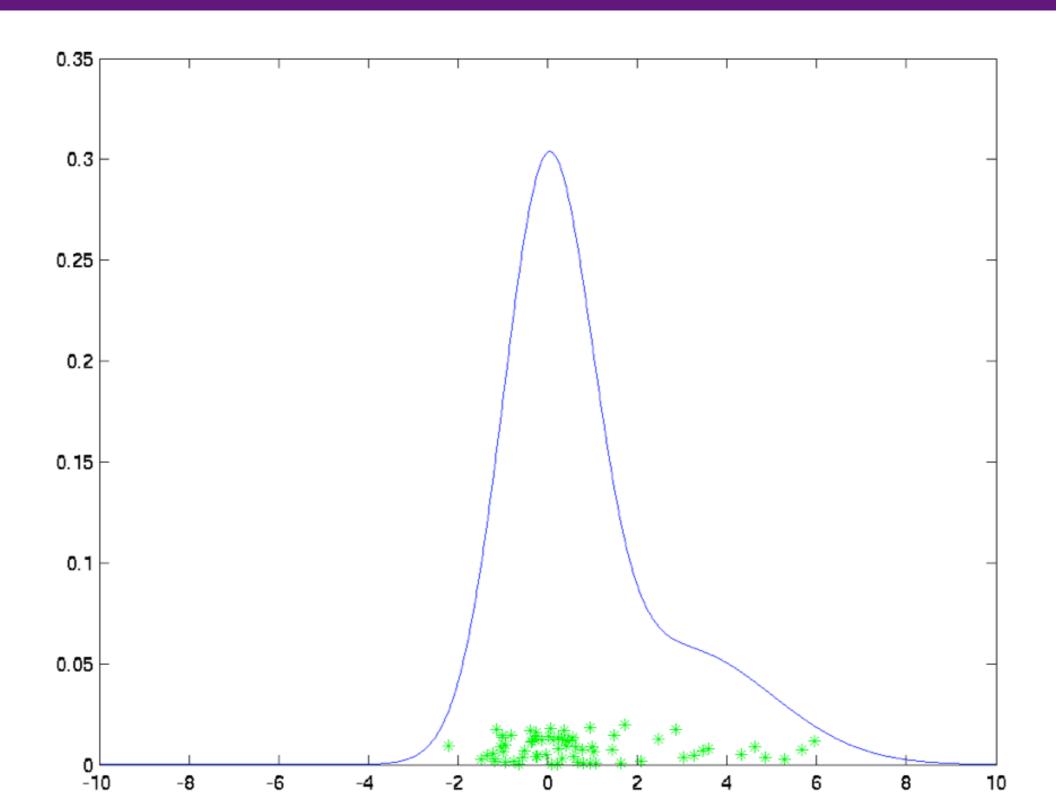


- Continuous domain = infinite number of bins
- Curse of dimensionality
  - 10 bins on [0, 1] is probably good
  - 10<sup>10</sup> bins on [0, 1]<sup>10</sup> requires high accuracy in estimate: probability mass per cell also decreases by 10<sup>10</sup>.

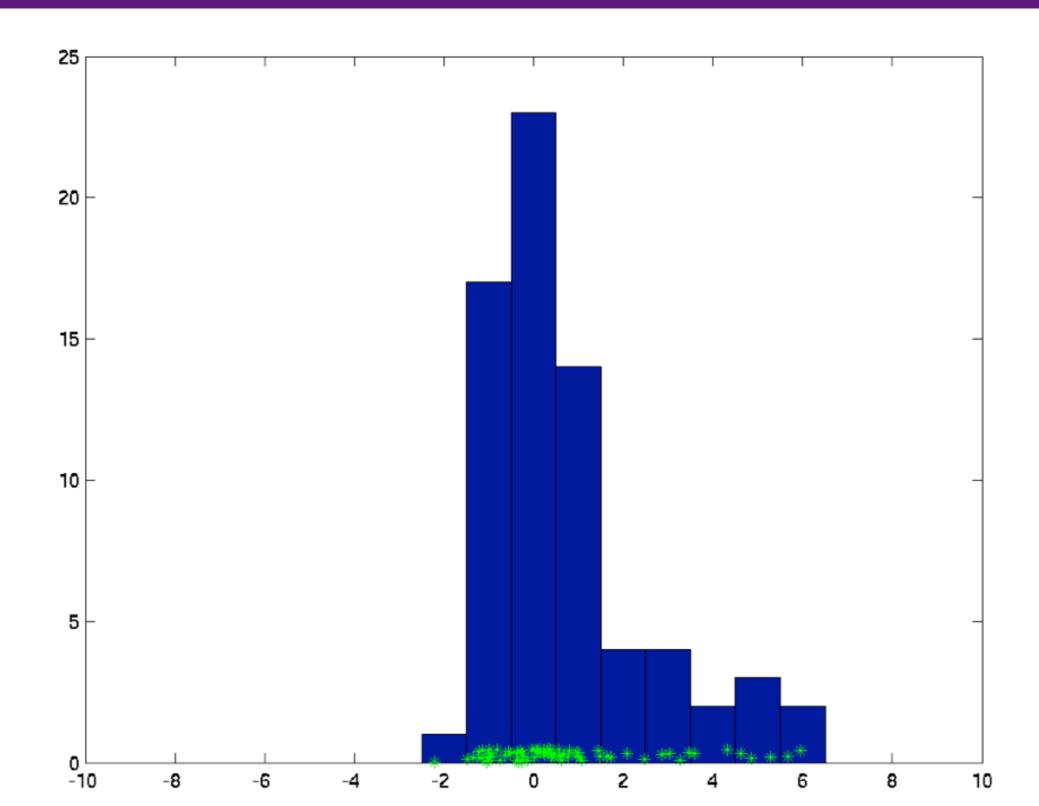
# Bin Counting



# Bin Counting



# Bin Counting



### Parzen Windows

Naive approach
 Use empirical density (delta distributions)

$$p_{\rm emp}(x) = \frac{1}{m} \sum_{i=1}^{m} \delta_{x_i}(x)$$

- This breaks if we see slightly different instances
- Kernel density estimate
   Smear out empirical density with a nonnegative smoothing kernel k<sub>x</sub>(x') satisfying

$$\int_{\mathcal{X}} k_x(x')dx' = 1 \text{ for all } x$$

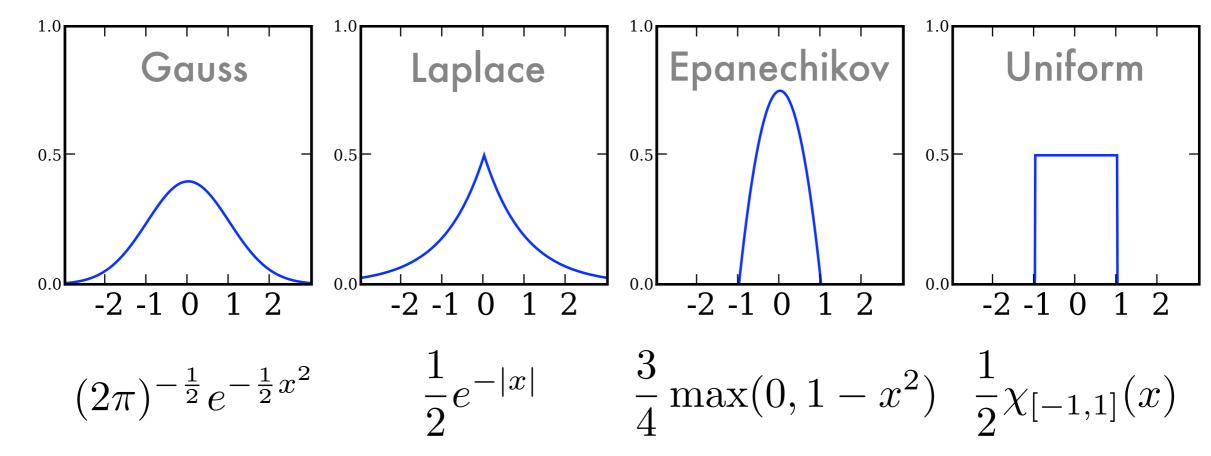
#### Parzen Windows

#### Density estimate

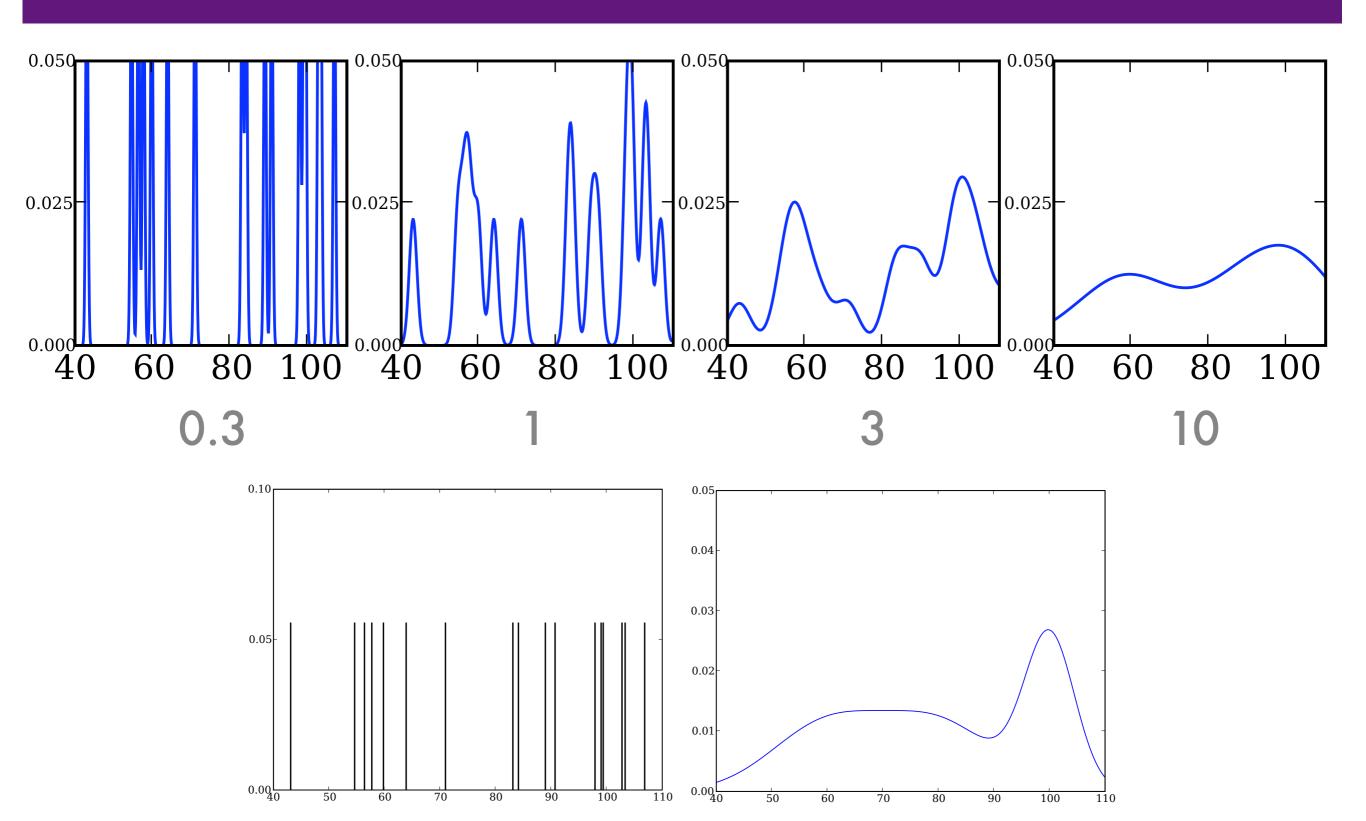
$$p_{\rm emp}(x) = \frac{1}{m} \sum_{i=1}^{m} \delta_{x_i}(x)$$

$$\hat{p}(x) = \frac{1}{m} \sum_{i=1}^{m} k_{x_i}(x)$$

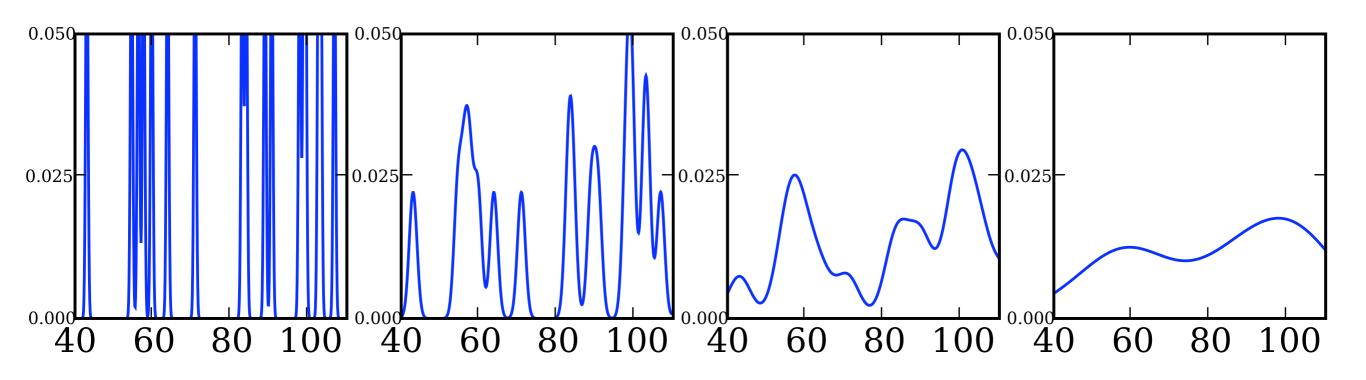
#### Smoothing kernels



### Size matters



#### Size matters

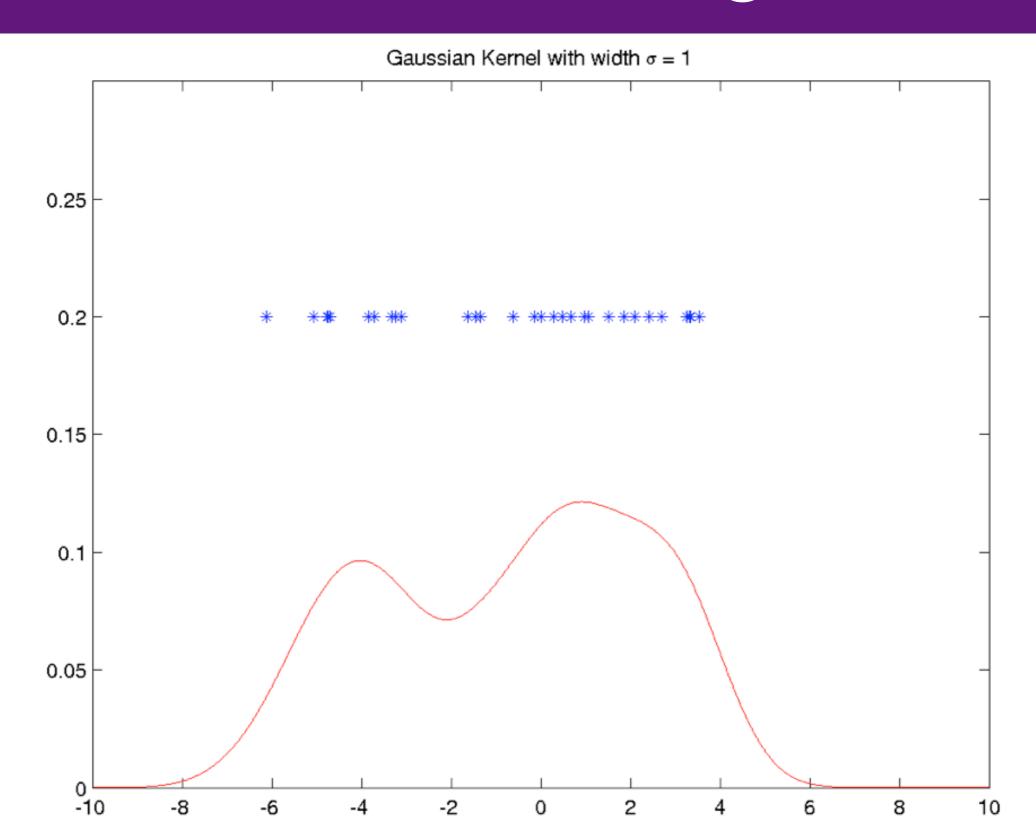


Kernel width

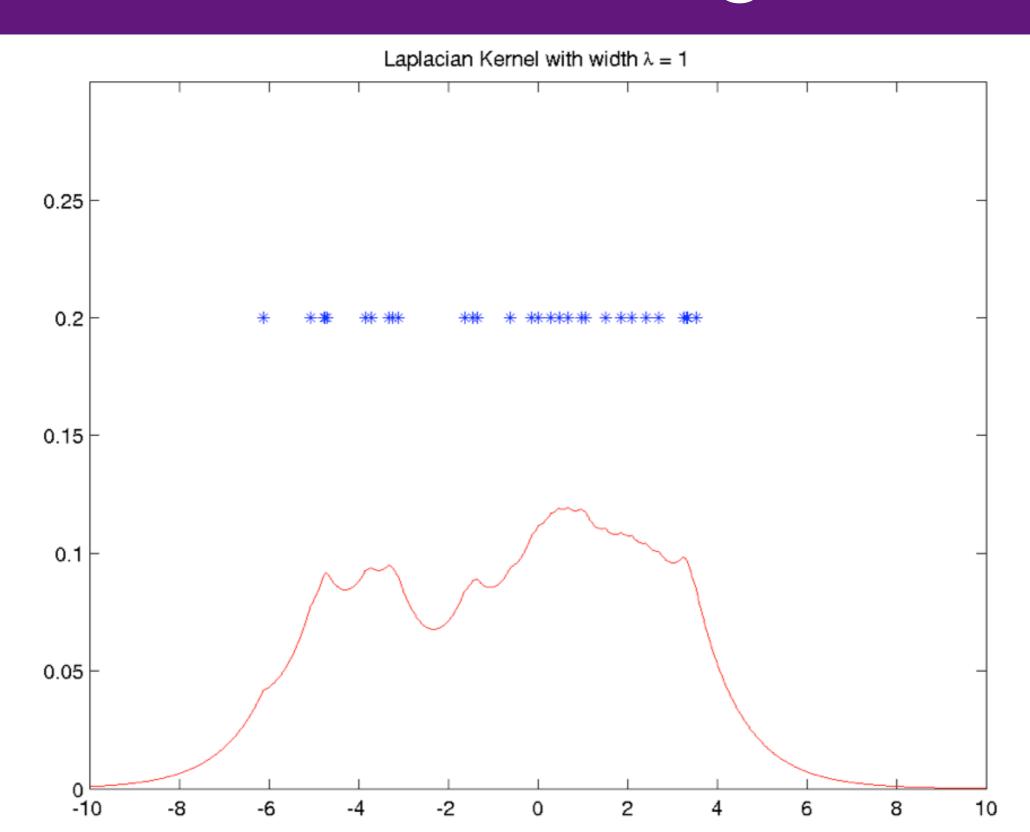
$$k_{x_i}(x) = r^{-d}h\left(\frac{x - x_i}{r}\right)$$

- Too narrow overfits
- Too wide smoothes with constant distribution
- How to choose?

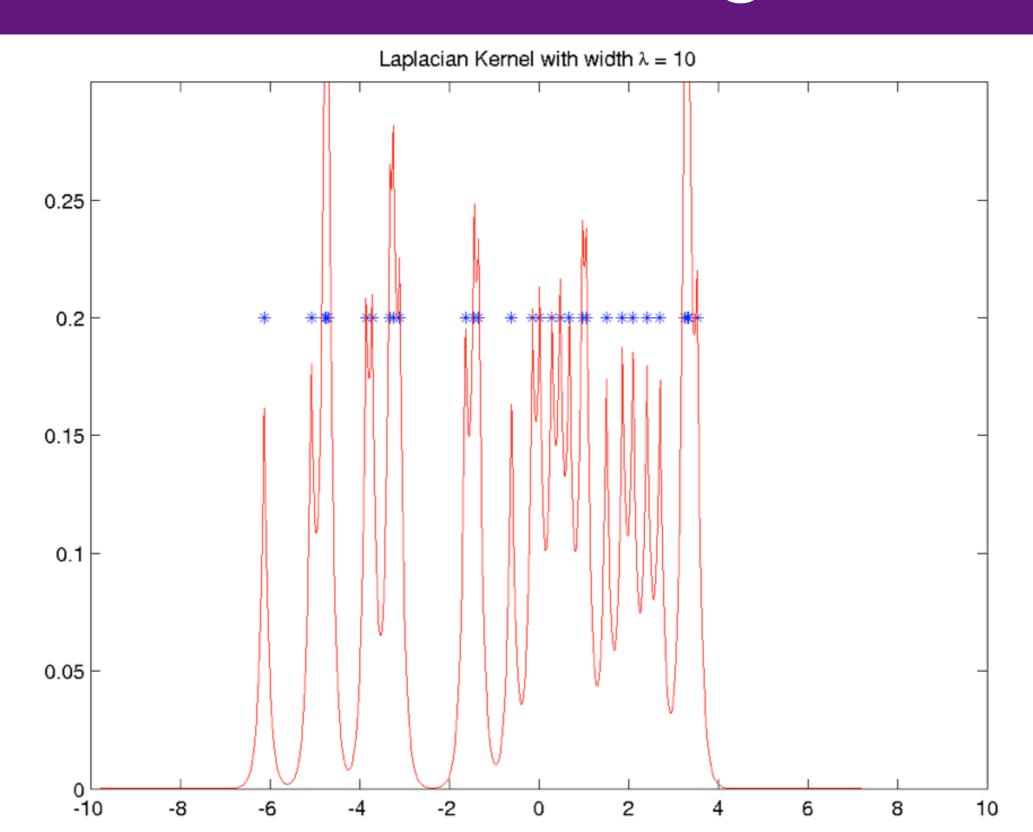
# Smoothing



# Smoothing



# Smoothing



# Capacity Control



# Capacity control

- Need automatic mechanism to select scale
- Overfitting
  - Maximum likelihood will lead to r=0 (smoothing kernels peak at instances)
  - This is (typically) a set of measure 0.
- Validation set
   Set aside data just for calibrating r
- Leave-one-out estimation
   Estimate likelihood using all but one instance
- Alternatives: use a prior on r; convergence analysis

# Capacity Control

#### Validation set

$$\log \hat{p}(X') = \sum_{x' \in X'} \log \hat{p}(x')$$

$$= \sum_{x' \in X'} \log \sum_{x \in X} k\left(\frac{x-x'}{r}\right) - |X'| \left[d\log r + \log|X|\right]$$

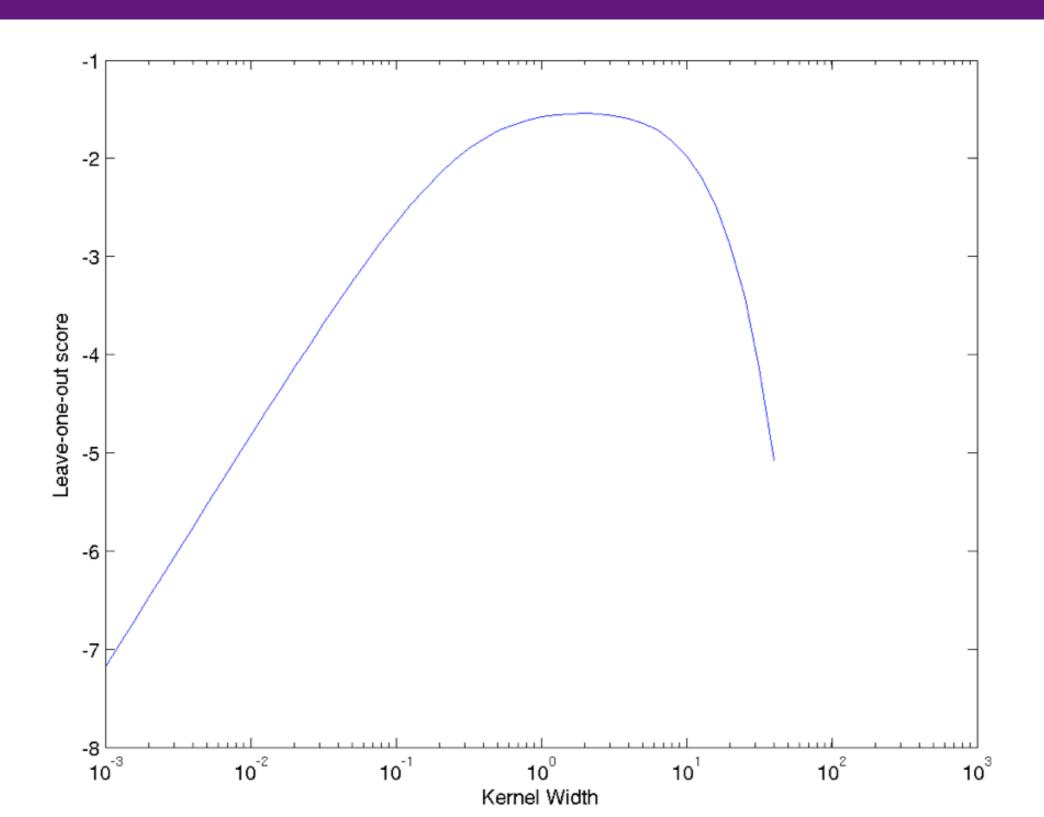
#### Leave-one-out crossvalidation

$$\hat{p}_{X\backslash\{x\}}(x) = \frac{1}{m-1} \sum_{x'\in X\backslash\{x\}} r^{-d}k\left(\frac{x'-x}{r}\right)$$

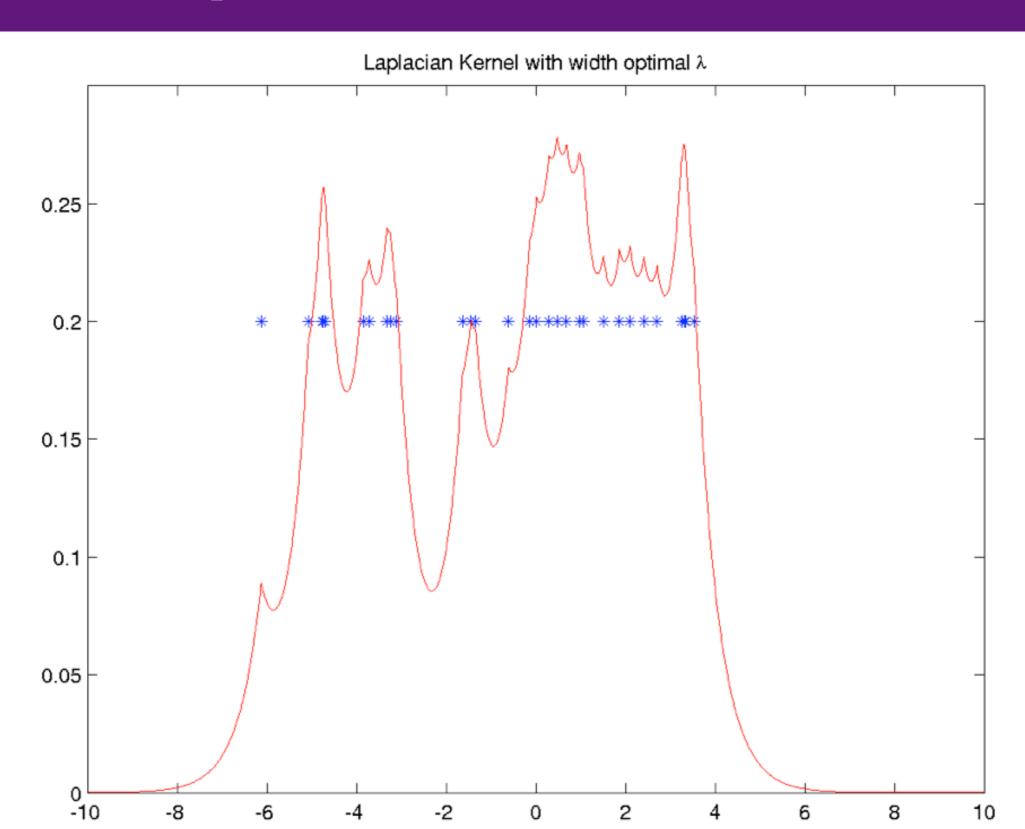
$$= \frac{m}{m-1} \left[\hat{p}(x) - m^{-1}r^{-d}k(0)\right]$$

$$\Longrightarrow \mathcal{L}[X] = m\log m/(m-1) + \sum_{x\in X} \log\left[\hat{p}(x) - m^{-1}r^{-d}k(0)\right]$$

### Leave-one out estimate



# Optimal estimate

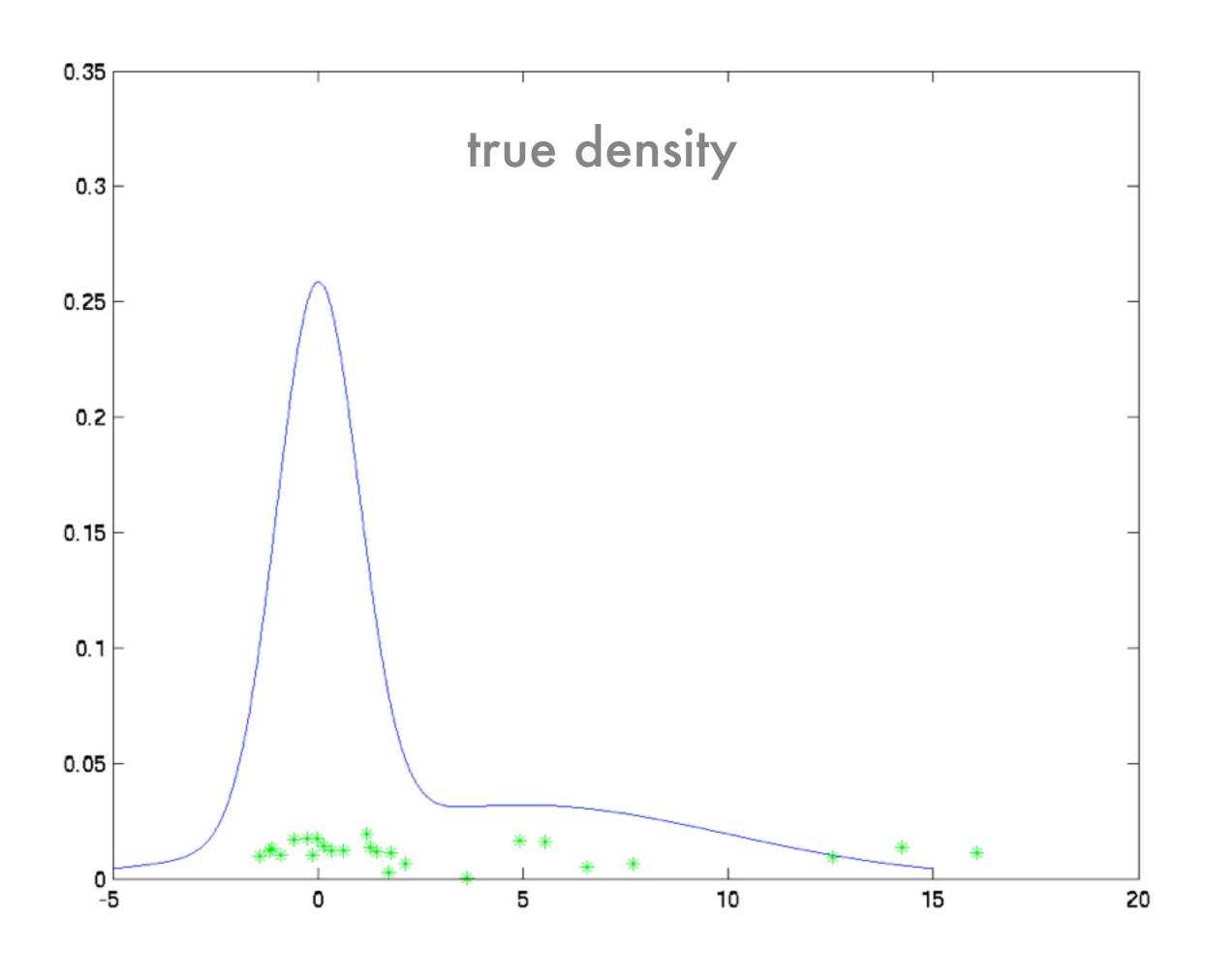


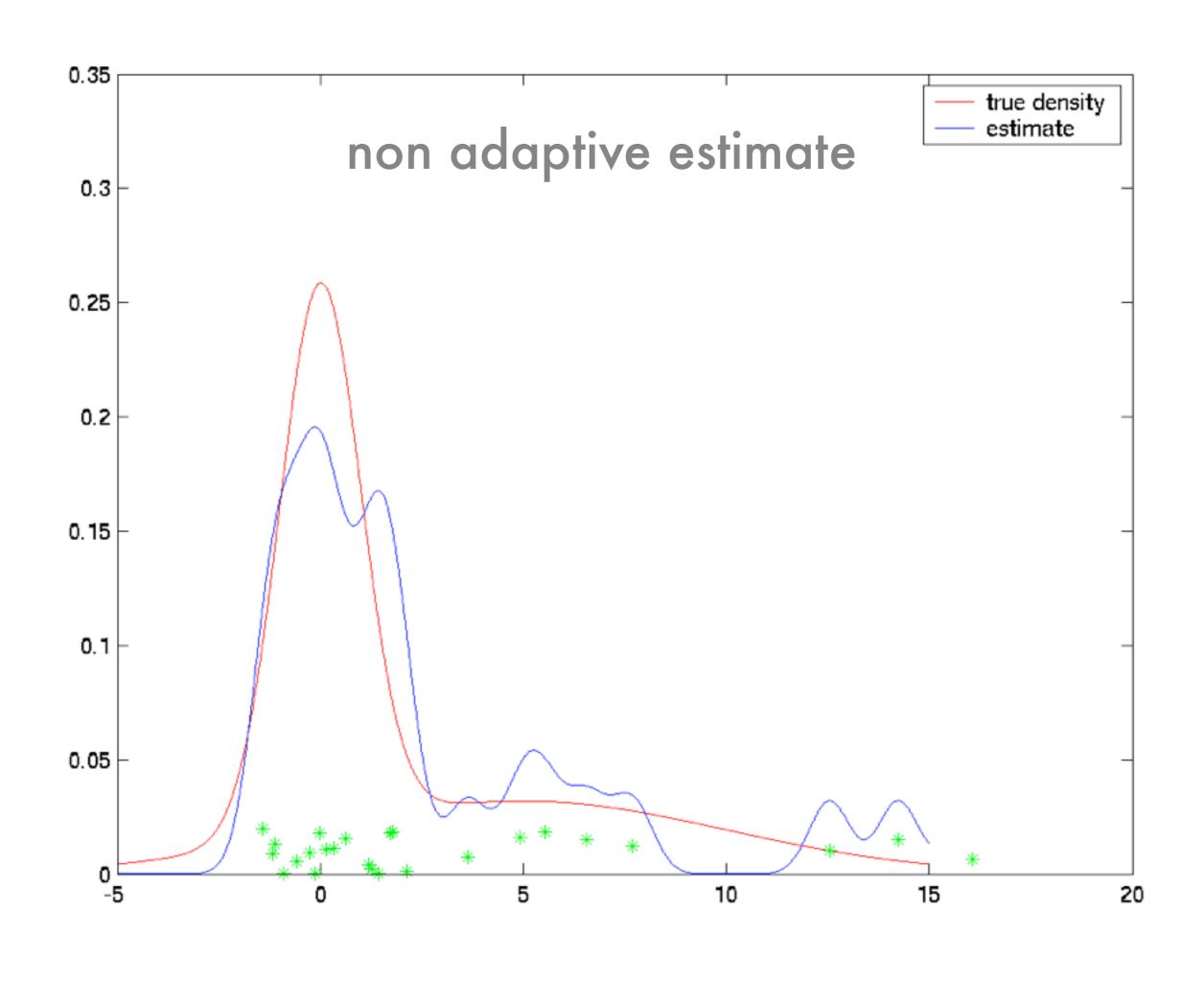
### Silverman's rule

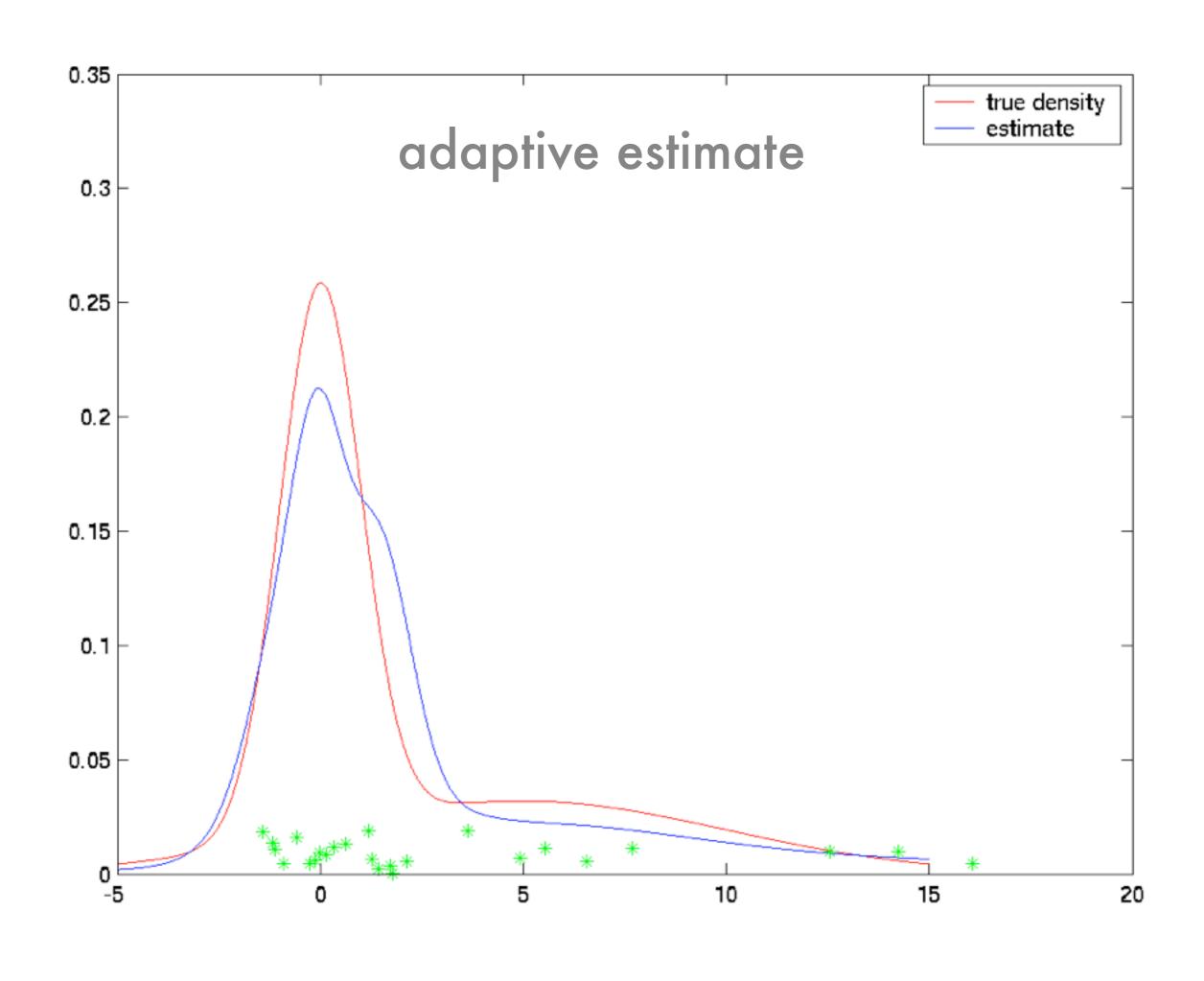
### Silverman's rule

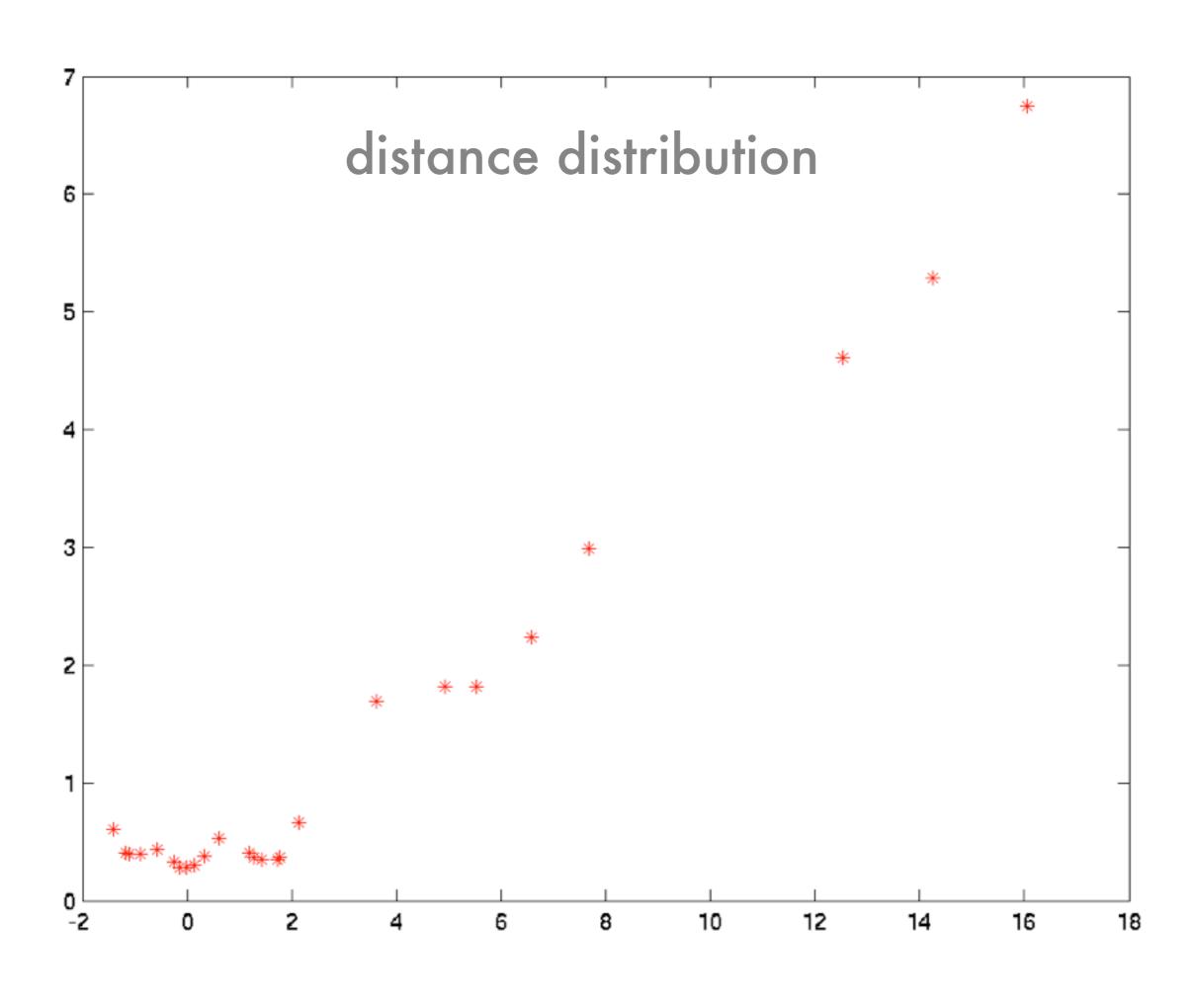
- Chicken and egg problem
  - Want wide kernel for low density region
  - Want narrow kernel where we have much data
  - Need density estimate to estimate density
- Simple hack
   Use average distance from k nearest neighbors

$$r_i = \frac{r}{k} \sum_{x \in \text{NN}(x_i, k)} ||x_i - x||$$









# Watson-Nadaraya estimator

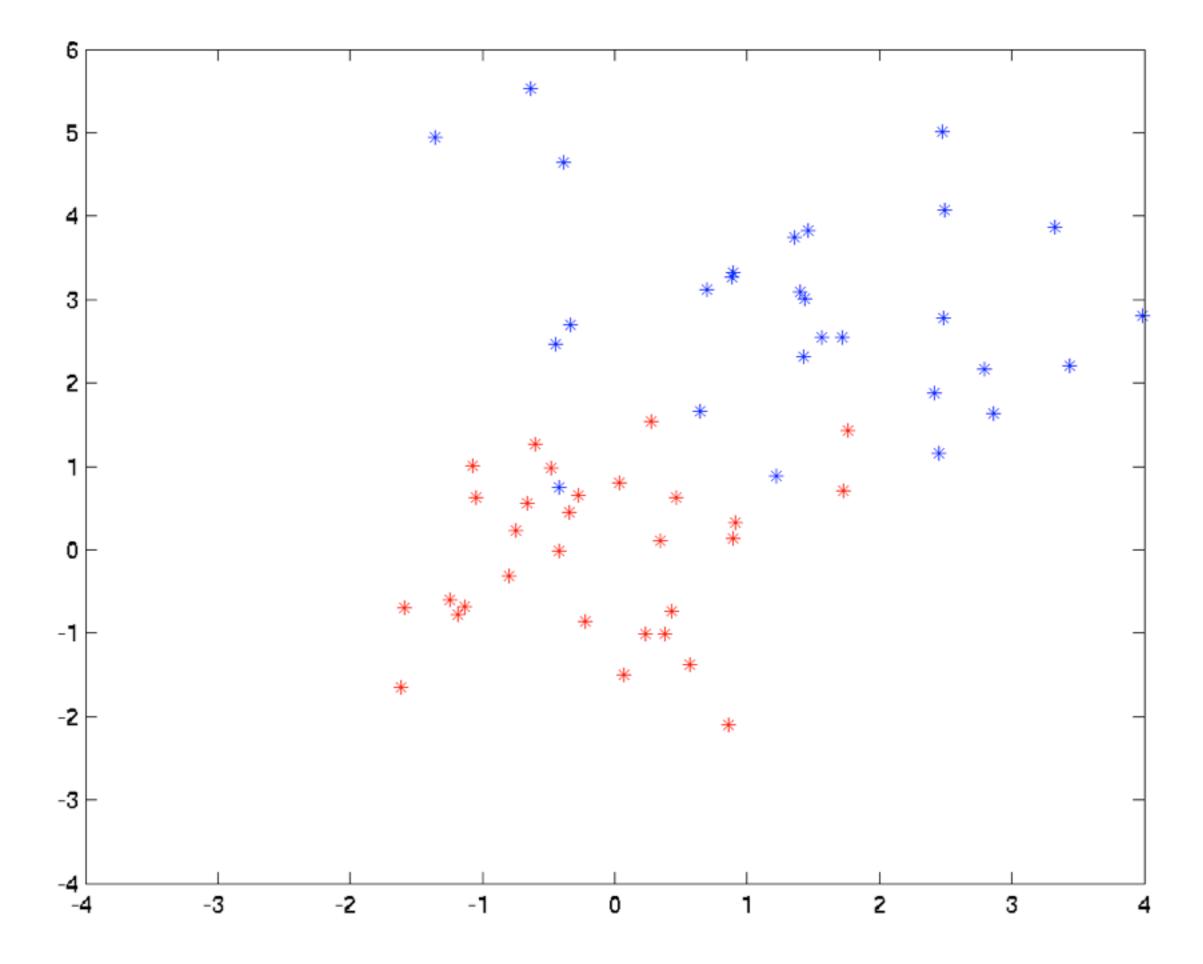
## Weighted smoother

- Problem
   Given pairs (x<sub>i</sub>, y<sub>i</sub>) estimate y | x for new x
- Idea
   Use distance weighted average of y<sub>i</sub>

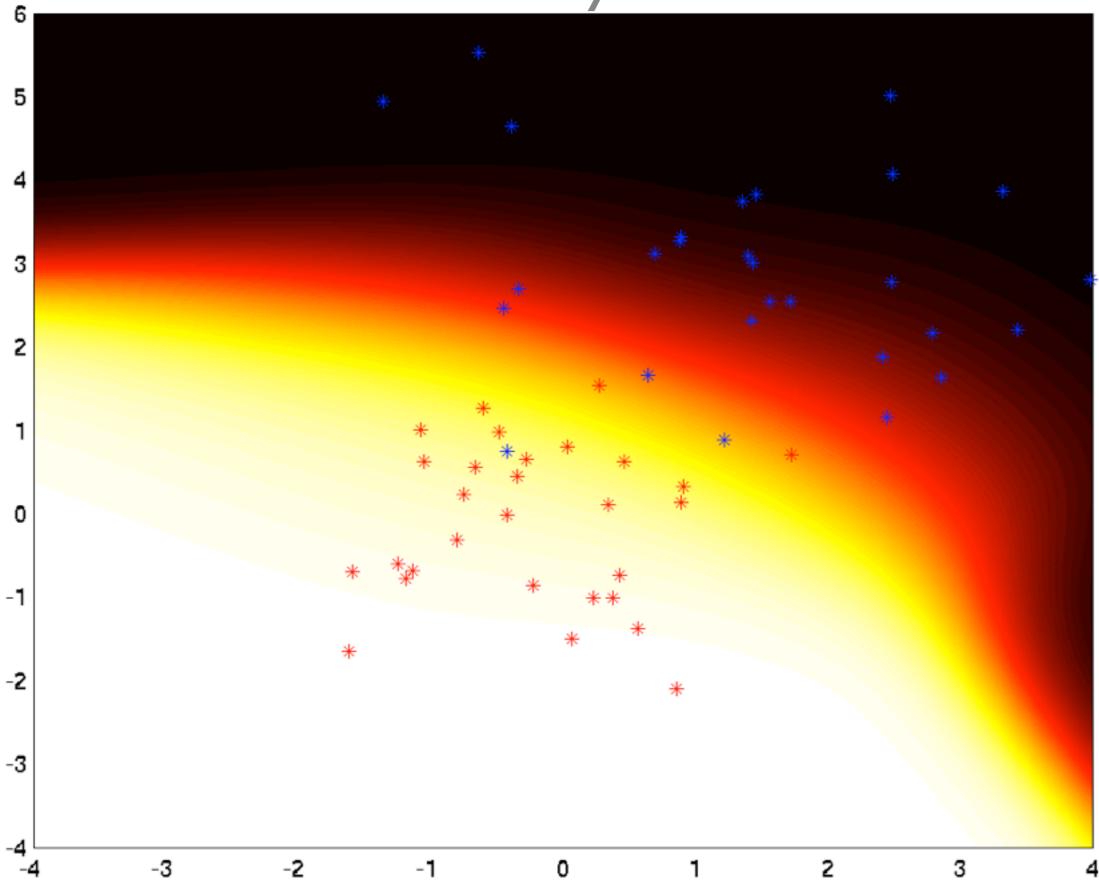
$$\hat{y}(x) = \sum_{i} y_i \frac{k_{x_i}(x)}{\sum_{j} k_{x_j}(x)} = \frac{\sum_{i} y_i k_{x_i}(x)}{\sum_{j} k_{x_j}(x)}$$

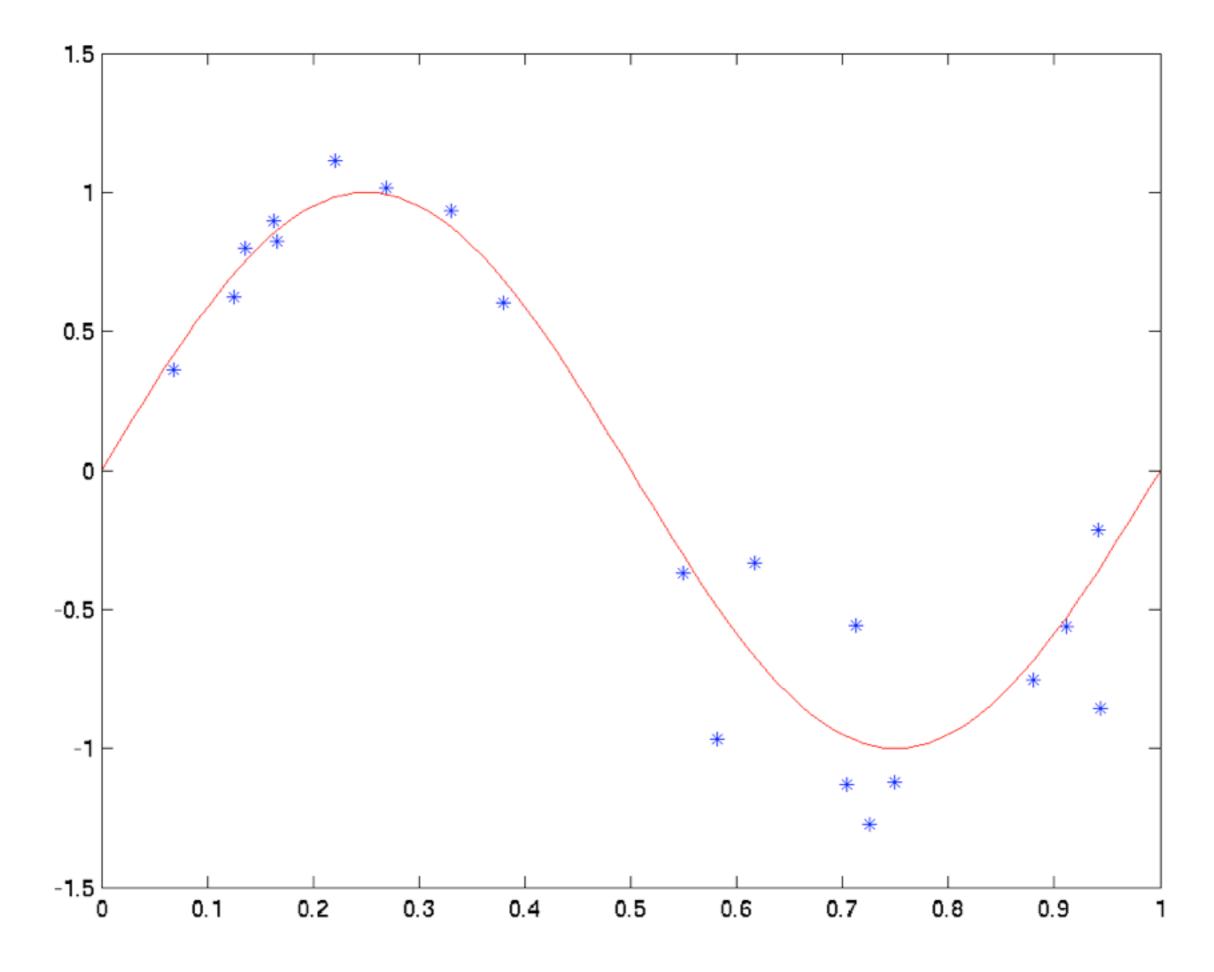
labels

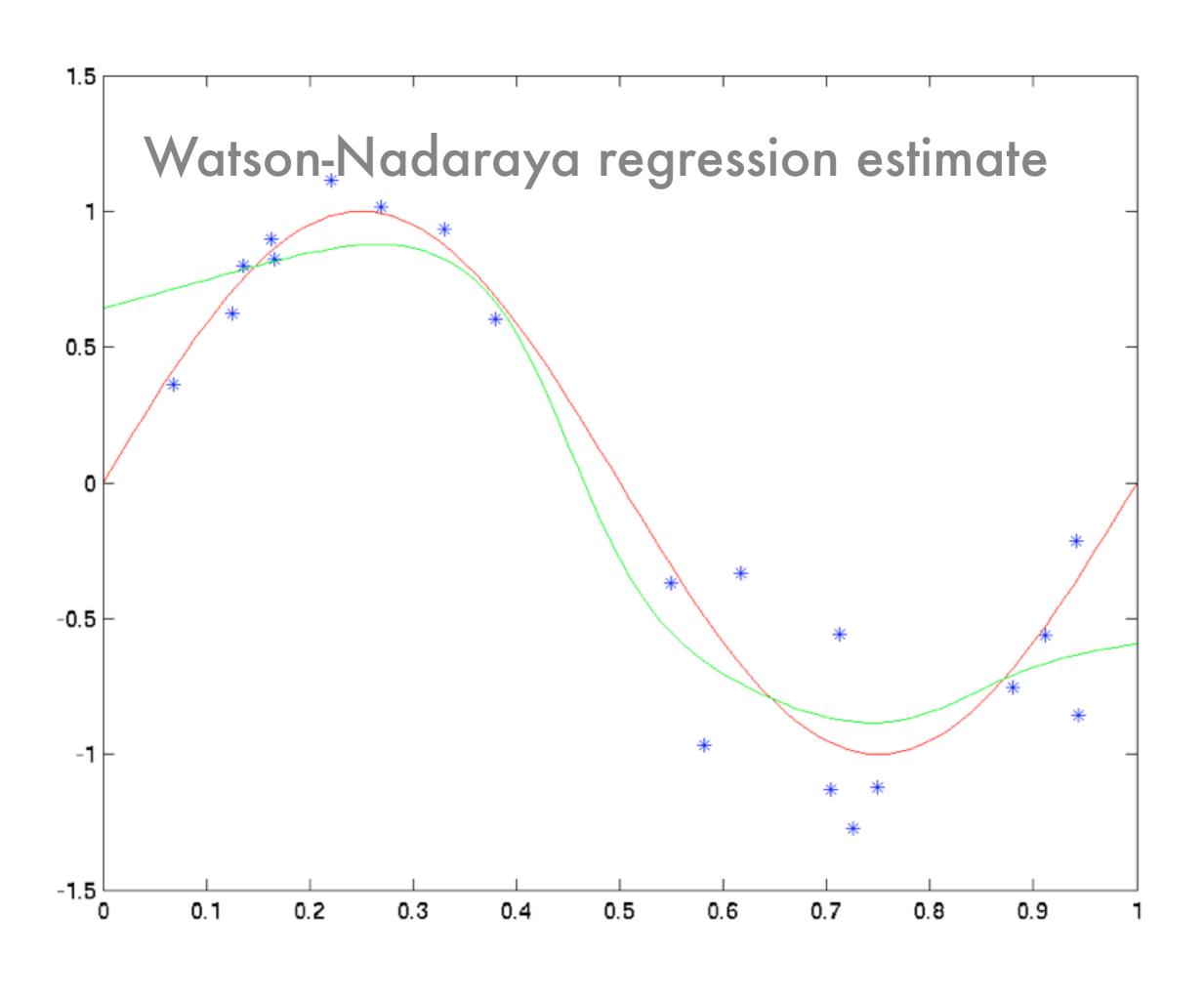
local weights



Watson-Nadaraya Classifier







## k-Nearest Neighbors

- Further simplification
  - Same weight for all nearest neighbors
  - Same number of neighbors everywhere
- Classification
   Use majority rule to estimate label
- Regression
   Use average for label



#### Density function

$$p(x; \theta) = \exp\left(\langle \phi(x), \theta \rangle - g(\theta)\right)$$
where  $g(\theta) = \log \sum_{x'} \exp\left(\langle \phi(x'), \theta \rangle\right)$ 

#### Density function

$$p(x; \theta) = \exp\left(\langle \phi(x), \theta \rangle - g(\theta)\right)$$
where  $g(\theta) = \log \sum_{x'} \exp\left(\langle \phi(x'), \theta \rangle\right)$ 

Log partition function generates cumulants

$$\partial_{\theta} g(\theta) = \mathbf{E} \left[ \phi(x) \right]$$
  
 $\partial_{\theta}^{2} g(\theta) = \operatorname{Var} \left[ \phi(x) \right]$ 

Density function

$$p(x; \theta) = \exp\left(\langle \phi(x), \theta \rangle - g(\theta)\right)$$
where  $g(\theta) = \log \sum_{x'} \exp\left(\langle \phi(x'), \theta \rangle\right)$ 

Log partition function generates cumulants

$$\partial_{\theta} g(\theta) = \mathbf{E} \left[ \phi(x) \right]$$
  
 $\partial_{\theta}^{2} g(\theta) = \operatorname{Var} \left[ \phi(x) \right]$ 

• g is convex (second derivative is p.s.d.)

## Examples

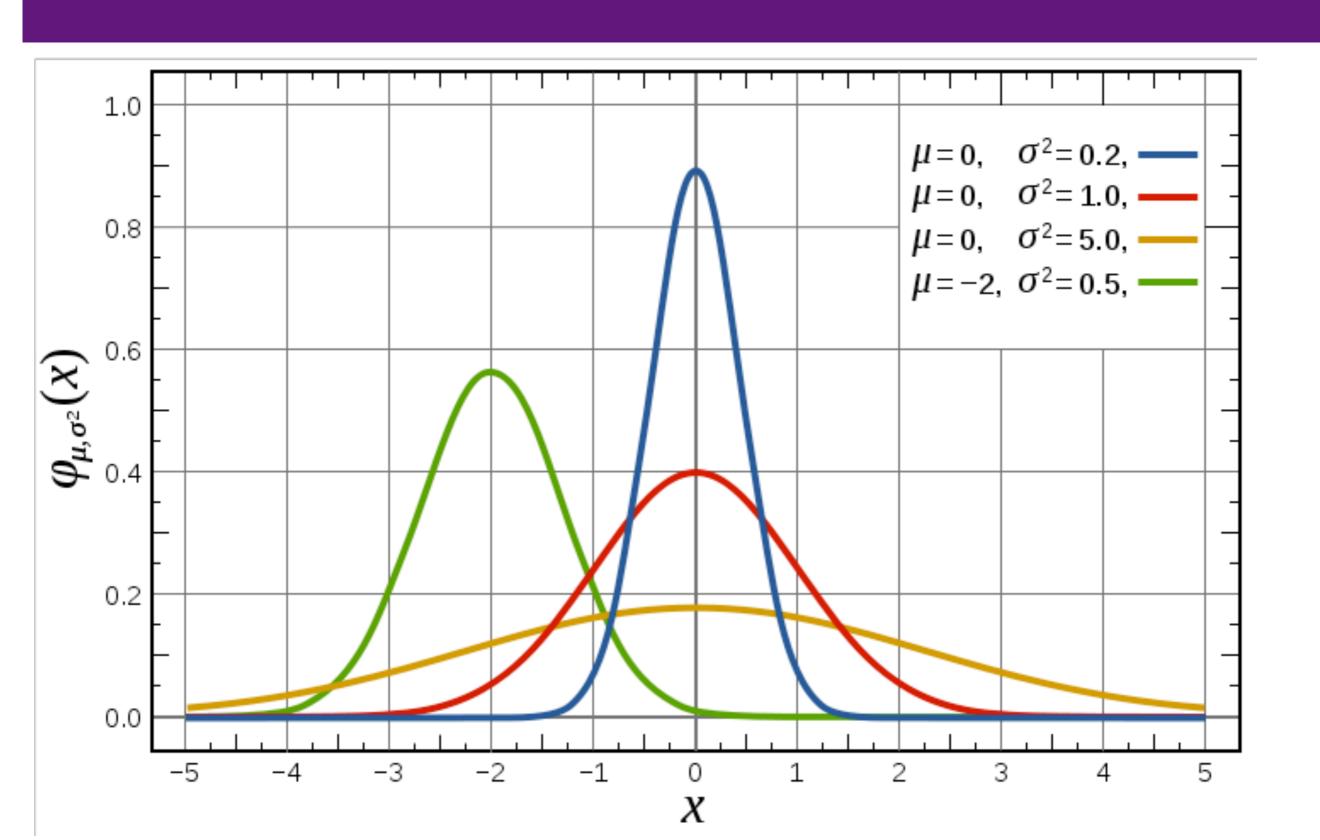
- Binomial Distribution
- Discrete Distribution
   (e<sub>x</sub> is unit vector for x)
- Gaussian
- Poisson (counting measure 1/x!)  $\phi(x) = x$
- Dirichlet, Beta, Gamma, Wishart, ...

$$\phi(x) = x$$

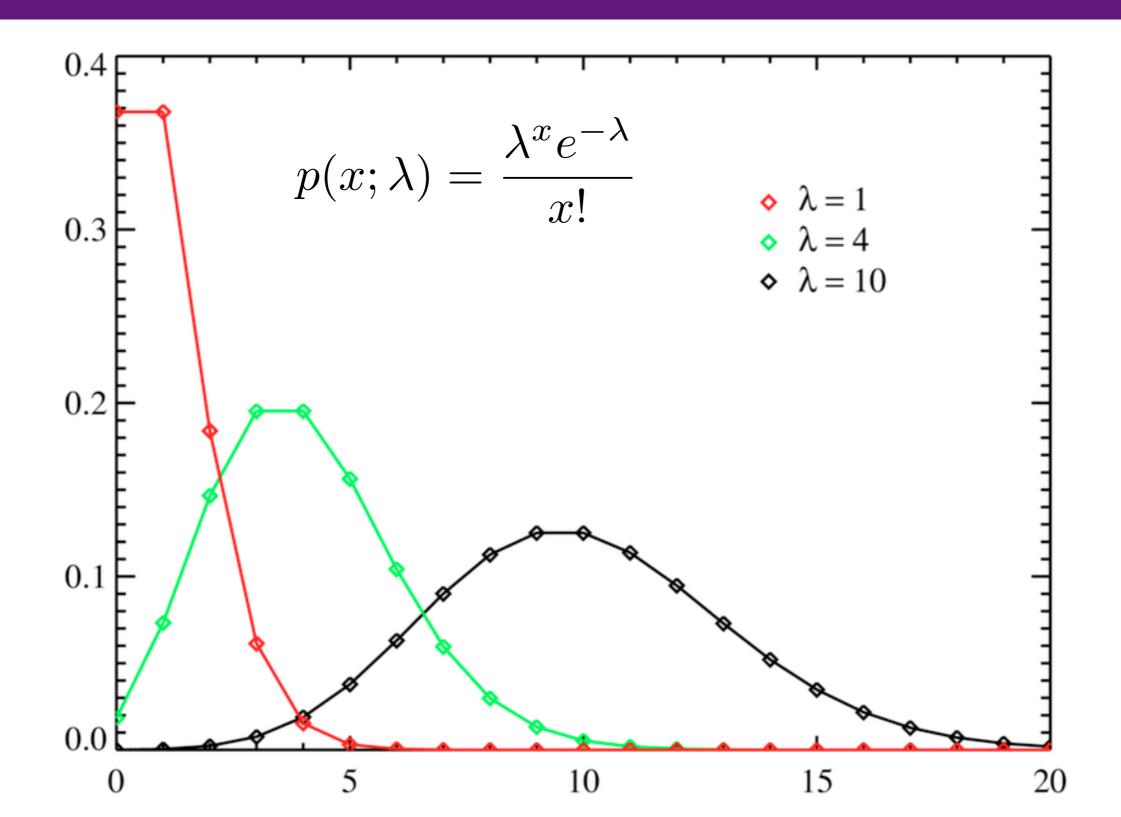
$$\phi(x) = e_x$$

$$\phi(x) = \left(x, \frac{1}{2}xx^{\top}\right)$$

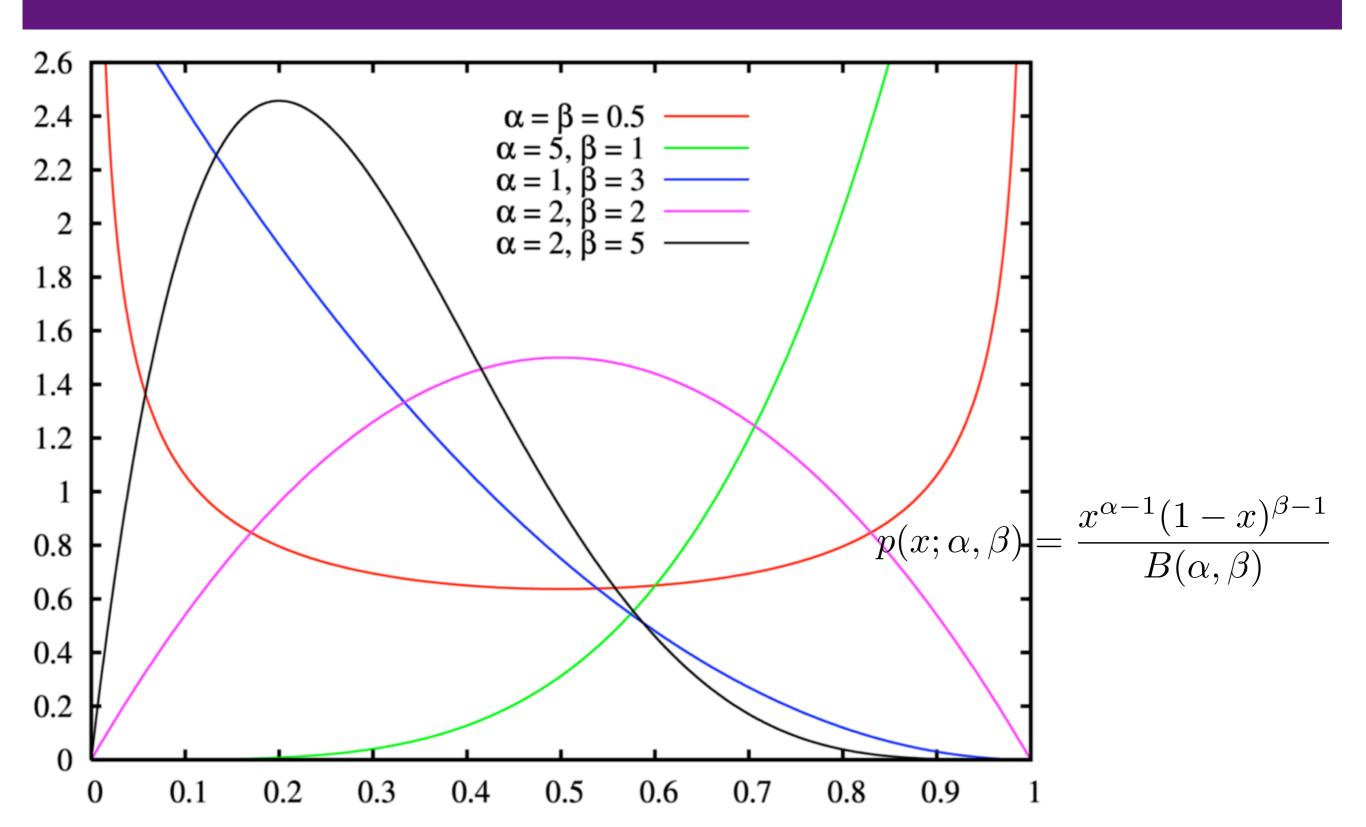
## Normal Distribution



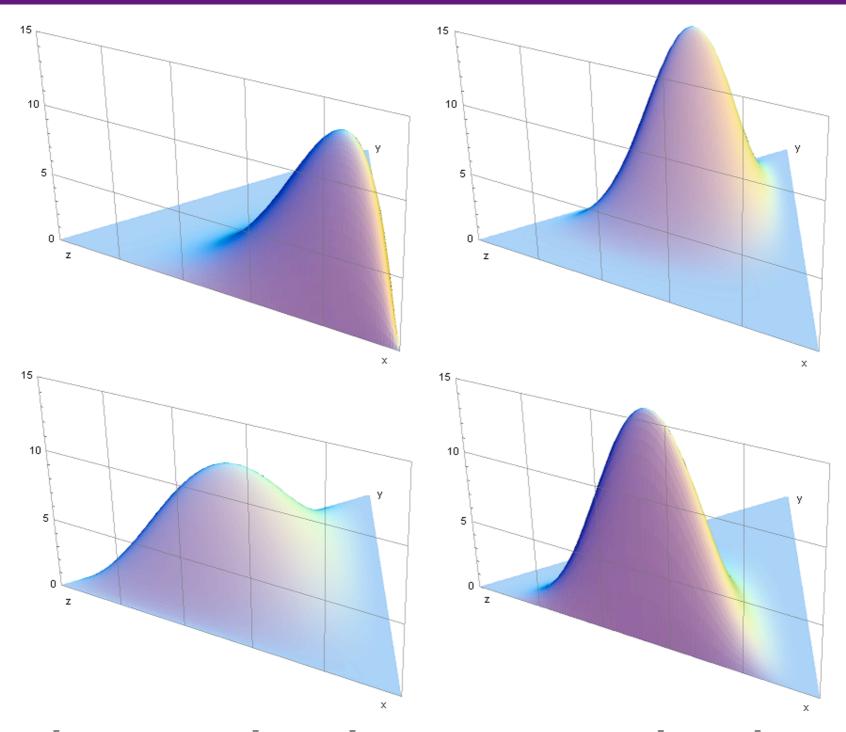
### Poisson Distribution



#### Beta Distribution



#### Dirichlet Distribution



... this is a distribution over distributions ...

### Maximum Likelihood

#### Maximum Likelihood

Negative log-likelihood

$$-\log p(X;\theta) = \sum_{i=1}^{n} g(\theta) - \langle \phi(x_i), \theta \rangle$$

#### Maximum Likelihood

Negative log-likelihood

$$-\log p(X;\theta) = \sum_{i=1}^{n} g(\theta) - \langle \phi(x_i), \theta \rangle$$

• Taking derivatives 
$$-\partial_{\theta}\log p(X;\theta) = m\left[\mathbf{E}[\phi(x)] - \frac{1}{m}\sum_{i=1}^{n}\phi(x_i)\right]$$

We pick the parameter such that the distribution matches the empirical average.

empirical

average

## Conjugate Priors

- Unless we have lots of data estimates are weak
- Usually we have an idea of what to expect

$$p(\theta|X) \propto p(X|\theta) \cdot p(\theta)$$

we might even have 'seen' such data before

Solution: add 'fake' observations

$$p(\theta) \propto p(X_{\text{fake}}|\theta) \text{ hence } p(\theta|X) \propto p(X|\theta)p(X_{\text{fake}}|\theta) = p(X \cup X_{\text{fake}}|\theta)$$

Inference (generalized Laplace smoothing)

$$\frac{1}{n}\sum_{i=1}^n\phi(x_i)\longrightarrow\frac{1}{n+m}\sum_{i=1}^n\phi(x_i)+\frac{m}{n+m}\mu_0$$
 fake count

## Example: Gaussian Estimation

- Sufficient statistics:  $x, x^2$
- Mean and variance given by

$$\mu = \mathbf{E}_x[x] \text{ and } \sigma^2 = \mathbf{E}_x[x^2] - \mathbf{E}_x^2[x]$$

Maximum Likelihood Estimate

$$\hat{\mu} = \frac{1}{n} \sum_{i=1}^{n} x_i \text{ and } \sigma^2 = \frac{1}{n} \sum_{i=1}^{n} x_i^2 - \hat{\mu}^2$$

Maximum a Posteriori Estimate

smoother

$$\hat{\mu} = \frac{1}{n+n_0} \sum_{i=1}^{n} x_i \text{ and } \sigma^2 = \frac{1}{n+n_0} \sum_{i=1}^{n} x_i^2 + \frac{n_0}{n+n_0} \mathbf{1} - \hat{\mu}^2$$

# Collapsing

Conjugate priors

$$p(\theta) \propto p(X_{\rm fake}|\theta)$$

Hence we know how to compute normalization

• Prediction  $p(x|X) = \int p(x|\theta)p(\theta|X)d\theta$ 

(Beta, binomial)
(Dirichlet, multinomial)
(Gamma, Poisson)
(Wishart, Gauss)

$$\propto \int p(x|\theta)p(X|\theta)p(X_{\rm fake}|\theta)d\theta$$
 
$$= \int p(\{x\} \cup X \cup X_{\rm fake}|\theta)d\theta$$
 look up closed form expansions

http://en.wikipedia.org/wiki/Exponential\_family

## Conjugate Prior in action

 $m_i = m \cdot [\mu_0]_i$ 

$$p(x=i) = \frac{n_i}{n} \longrightarrow p(x=i) = \frac{n_i + m_i}{n+m}$$

Outcome	1	2	3	4	5	6
Counts	3	6	2	1	4	4
$\overline{\mathrm{MLE}}$	0.15	0.30	0.10	0.05	0.20	0.20
MAP $(m_0 = 6)$	0.15	0.27	0.12	0.08	0.19	0.19
MAP $(m_0 = 100)$	0.16	0.19	0.16	0.15	0.17	0.17

## Conjugate Prior in action

 $m_i = m \cdot [\mu_0]_i$ 

#### Discrete Distribution

$$p(x=i) = \frac{n_i}{n} \longrightarrow p(x=i) = \frac{n_i + m_i}{n+m}$$
   
 • Tossing a dice

Outcome	1		3	4	5	6
Counts	3	6	2	1	4	$\boxed{4}$
MLE	0.15	0.30	0.10	0.05	0.20	0.20
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## Conjugate Prior in action

$$m_i = m \cdot [\mu_0]_i$$

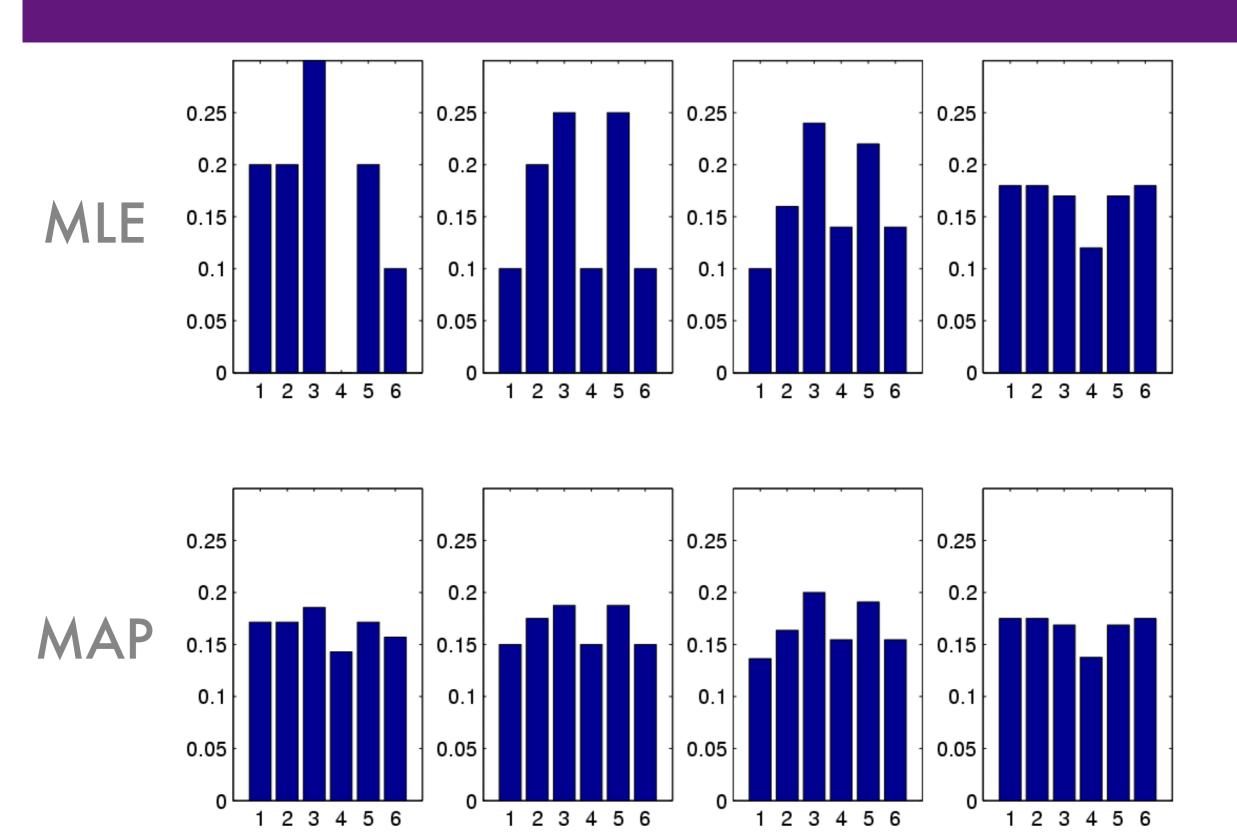
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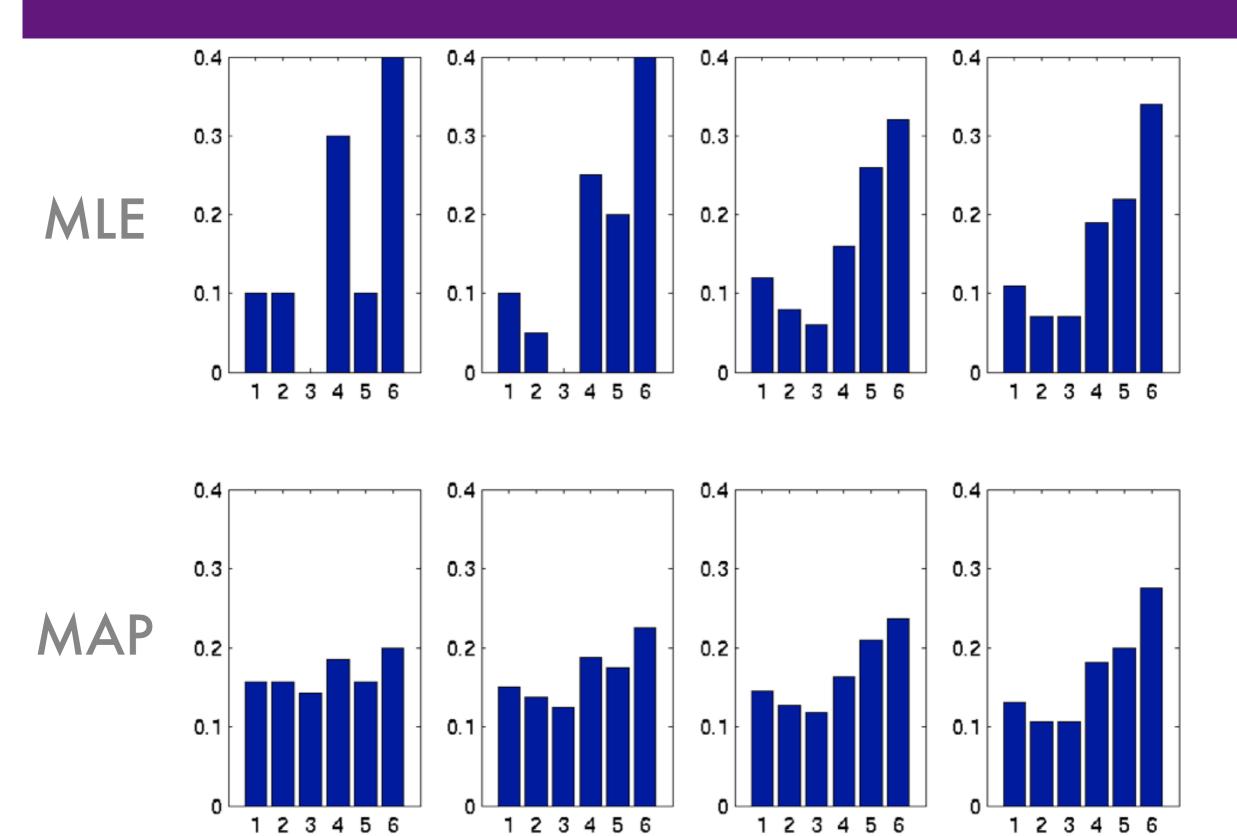
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 Rule of thumb need 10 data points (or prior) per parameter

#### Honest dice



### Tainted dice



## Priors (part deux)

Parameter smoothing

$$p(\theta) \propto \exp(-\lambda \|\theta\|_1) \text{ or } p(\theta) \propto \exp(-\lambda \|\theta\|_2^2)$$

Posterior

$$p(\theta|x) \propto \prod_{i=1}^{m} p(x_i|\theta)p(\theta)$$

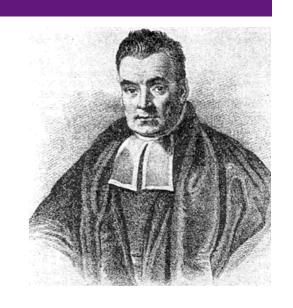
$$\propto \exp\left(\sum_{i=1}^{m} \langle \phi(x_i), \theta \rangle - mg(\theta) - \frac{1}{2\sigma^2} \|\theta\|_2^2\right)$$

Convex optimization problem (MAP estimation)

$$\underset{\theta}{\text{minimize }} g(\theta) - \left\langle \frac{1}{m} \sum_{i=1}^{m} \phi(x_i), \theta \right\rangle + \frac{1}{2m\sigma^2} \|\theta\|_2^2$$

#### Statistics

- Probabilities
  - Bayes rule, Dependence, independence, conditional probabilities
  - Priors, Naive Bayes classifier
  - Tail bounds
    - Chernoff, Hoeffding, Chebyshev, Gaussian
    - A/B testing
- Kernel density estimation
  - Parzen windows, Nearest neighbors,
     Watson-Nadaraya estimator
- Exponential families
  - Gaussian, multinomial, Poisson
  - Conjugate distributions and smoothing, integrating out



Peninsula Grill

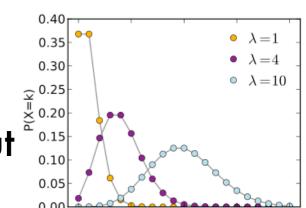


Come check out our new menu specials at your favorite city diner



Peninsula Grill

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## Further reading

Manuscript (book chapters 1 and 2)
 http://alex.smola.org/teaching/berkeley2012/slides/chapter1\_2.pdf